

Embedded Systems: Processor Classes

- General Purpose - high performance processors
 - Pentium, SPARC, ...
 - Used for general purpose software
 - General under control of an Operating System (Windows., UNIX..)
 - Workstation, PC's
- Embedded processors and special cores
 - 486SX, ARM, Hitachi SH7000, NEC V800, Motorola PowerPc, ...
 - Single program
 - Lightweight , often under realtime OS (Windows-CE, QNX, OS-9, PSOS, ...)
 - Support for -light- digital signal processing
 - Consumer electronic, cellular phones
- Microcontrollers
 - 8-, 16-, (32) Bit, application specific chipsets with various features
 - A/D, D/A, Input, Output, Timer, Interfaces,
 - Cost sensitive
 - Highest volume processors for automotive, washing machines,
- General purpose DSPs and core design in Gate-Arrays

Embedded DSP: Introduction

- **Digital Signal Processing**: application of mathematical operations to digital signals from the real world.
- Signals are represented digitally as a **sequence** of samples
- Samples obtained from physical signals via **transducers** (e.g. sensors) and **analog-to-digital converters** (ADC)
- Back-conversion to physical signals by **digital-to-analog converters** (DAC)
- Digital Signal Processors (DSP) is a electronic device that processes digital signals

Embedded DSP: Common Applications

- Applications
 - Scientific data processing
 - Communication
 - Audio and video processing
 - Graphics
 - Navigation
 - Control - robotics, guidance, machine vision
- Common signal processing algorithms
 - Filtering
 - Transformations (time \leftrightarrow frequency)
 - Convolution
 - Correlation
 - Matrix calculations

Embedded DSP: Target for DSPs

- DSP task demands:
 - Fast repetition of numeric computations
 - High memory and I/O-bandwidth to move data to the computation units
 - Deterministic reaction to signals from the physical world
 - Real-time processing capabilities
- DSPs must therefore perform these tasks efficiently while minimizing:
 - Power consumption
 - Memory usage for a given application
 - Development time
 - Cost
- DSPs need an excellent software support in order to have the full benefit of the hardware technology

Embedded DSP: DSP vs. General Purpose MPU

- DSPs generally run one program or a small, fixed set of programs
 - **No operating system** is generally necessary
 - Hence small special **OSes** are much simpler, there is no **virtual memory** (disk, ..) or protection
- DSPs often work in **hard real-time** applications
 - Necessary to perform the signal processing in a **fixed period of time**
 - Must be able to **react within a fixed time** to anything and everything that could happen
 - Interrupt handling: the interrupts or exceptions reduce the time for computation
 - Memory access and -latency: How to get data to the ALU/MAC and again out of the DSP ?
- DSPs often process an infinite **continuous** data sequence according to time slots from the environment (e.g. audio processing with 48 KHz sampling rate)

Embedded DSP: DSP vs. General Purpose MPU

- View to the market shows that in terms of dollar, the biggest markets for DSP processors today are:
 - Digital cellular telephony systems
 - Modems
 - Pagers and Bluetooth systems (wireless)
 - Disk drive servo and general drive controls
- The engineers demand
 - Good overall performance
 - Low cost
 - Energie efficiency
 - Support for hardware
 - Efficient support for software development
 - Support by GUI based tools (Mathlab, Labview, DSPWorks,)
 - Algorithm support by DSP-libraries

(Application specific choose of the **best matched processor**)

Embedded DSP: General Features of DSPs

- Single cycle CPUs (key of fast computation,..)
- SISD and SIMD structures (multiple computation units,..)
- Specialized instruction set (RISC) (efficient instructions,..)
- Integer and/or floating point units (Scaling, high signal dynamic,..)
- Specialized computation modules (Fast, parallel with file registers)
(ALU/MAC/SHIFTER)
- Data path configured for DSP applications (FIR, IIR: coefficients and data,..)
- Multiple memory banks and buses, cache (Parallel acces,..)
- Specialized addressing modes (Bit-reverse for FFT,..)
- Specialized execution control (Fast context switching,..)
- Fast interrupt response (Real time)
- Specialized peripherals for input/output of data (Multiport, DMA,..)
- Specialized units for multiprocessing (Fast LINK structures,..)

Embedded DSP: Addressing

- Standard DSPs have standard addressing modes (immediate, register indirect, direct)
- Additional performance by complex addressing modes
 - Want to keep MAC datapath performing computation
 - Offload addressing to somewhere else outside the chip
- Modes
 - Autoincrement/autodecrement (post/pre) before/after generating the address
 - Circular
 - Data producer/consumer functions with infinite I/O
 - Producer writes after tail pointer in autoincrement mode
 - Consumer reads from head pointer in autoincrement mode
 - The pointer wrap around by a special modulo arithmetic
 - Bit-reversed (e.g. usefull for FFT-Transform)
 - The pointers increment in bit-reversed order instead of normal incrementing

Embedded DSP: Multiplication

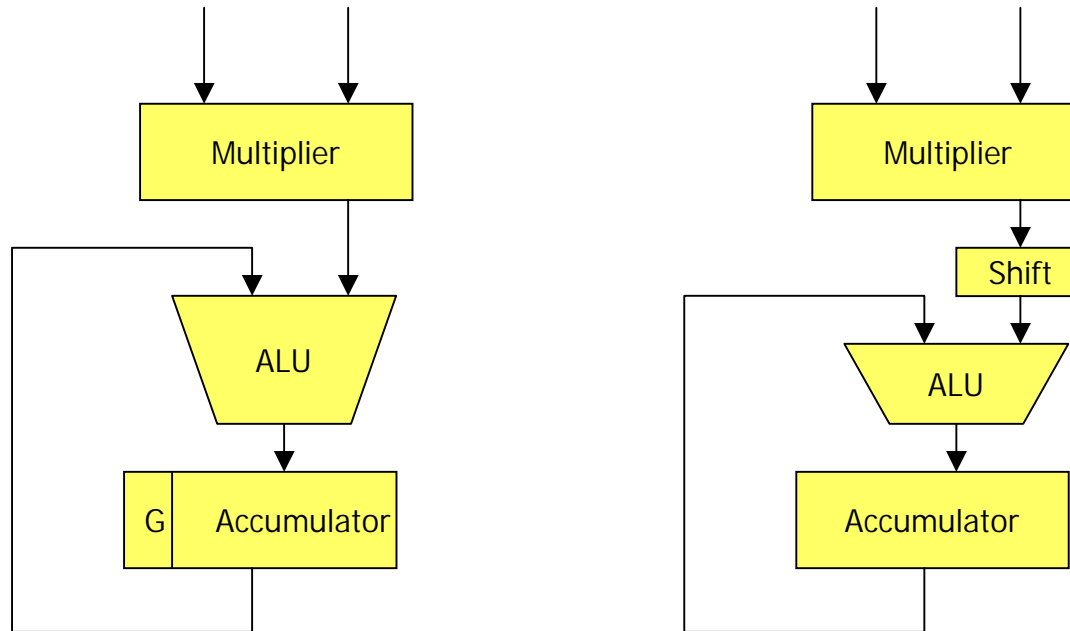
- One cycle multiplication
 - The most instruction in real signal algorithms are multiplies
 - A fast multiplier is a big demand
- One cycle multiply-accumulate (MAC)
 - $y[n] = y[n] + n * x[i]$ --> common operation in filtering (see last lecture !!)
- N-bit multiply gives an 2n-bit product
 - Number of bits for the MAC ?
- Need to round the results to minimize the bias
- Integer and/or floating point DSP ?

Embedded DSP: DSP Arithmetic

- Working with real signal - want real numbers
- DSPs support 16, 24, 32, 48, 64 (80) bit width
- Where is the decimal point ?
 - Programmer interprets the data in accordance to the task
 - Programmers will scale the data inside the function
 - Sometime a shift operation can prevent overflow
- DSPs work with data that was originally analog out of the real world
- Saturation arithmetic is necessary
 - Most DSP algorithms depend on special saturation arithmetic

Embedded DSP: Accumulate

- Demand: No overflow and no scale !
- Solution 1: Guard bits - make the accumulator wider than the product
- Solution 2: Scale by shifting - before adding to the accumulator



Embedded DSP: Memory

- DSP Algorithms are data-intensive
- Demand for a high data-throughput
- Minimize instruction bandwidth
 - Special zero overhead loops
 - Loop buffers
 - Circular buffer
 - Bit-Reverse technology
 - One key is a cache memory for small functions
- What about data bandwidth ?
 - Multiport Register File (e.g. MAC requires two new input values per cycle !)
 - Multiple data banks for program and data
 - Fast internal memory; but extern memory ?
 - What about program and/or data cache ?

Embedded DSP: Why not a CISC processor ?

- DSPs takes their power by doing a lot of things in parallel
 - MAC
 - Complicated addressing
 - Special I/O-controller hardware for data management
- Why not use a superscalar processor architecture ?

Superscalar is not predictable

- **Very Long Instruction Words (VLIW) is !**
- **Example: SHARC**

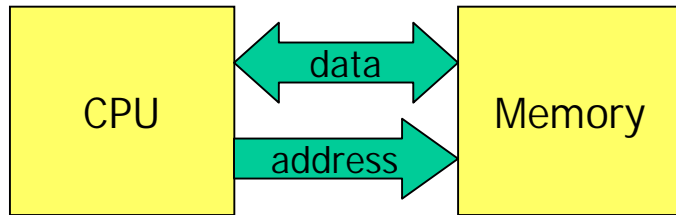
Embedded DSP: DSP versus GP Processors

- Specialized for fast computation on signals
- MIPS/MFLOPS is only MAC speed
- Performance of algorithm is the target
- Really fast only by special program techniques and a deep understanding of the hardware !

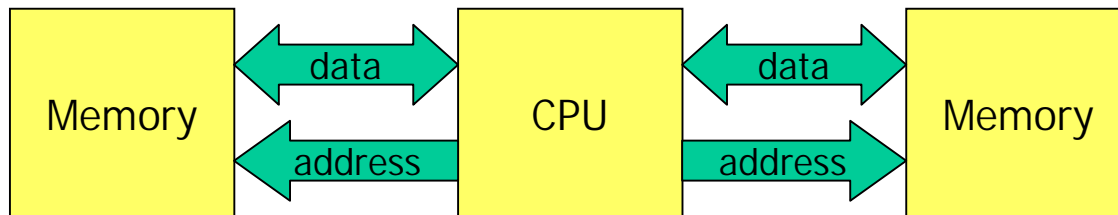
Are DSPs obsolete ?

- Additional performance improvements of general microprocessors and powerful media extensions increase the performance of GP processors to approach DSP performance.
- Can GP processors do the tasks of DSP ?

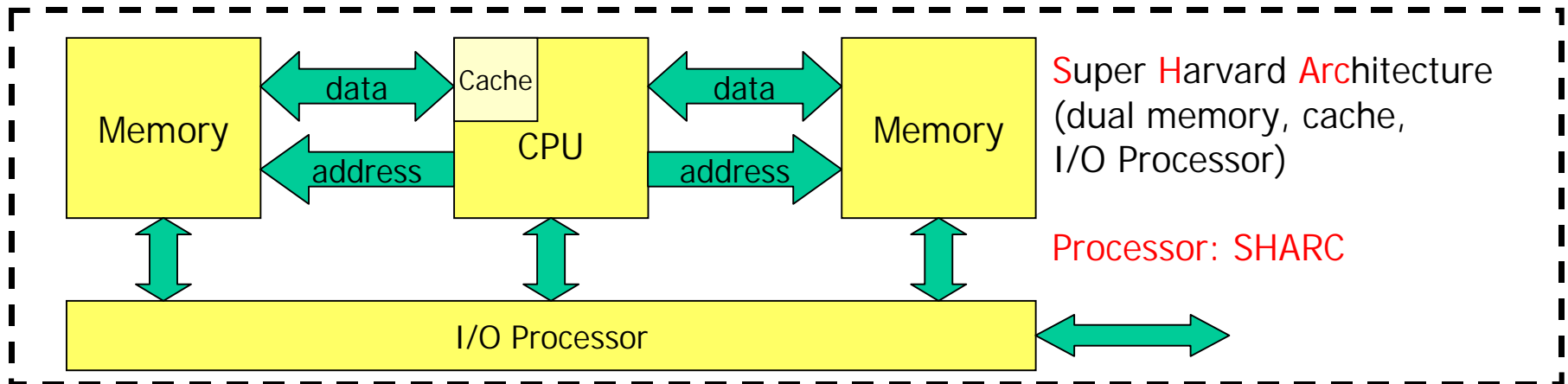
Embedded DSP: General Architectures



Von Neumann
(single memory)



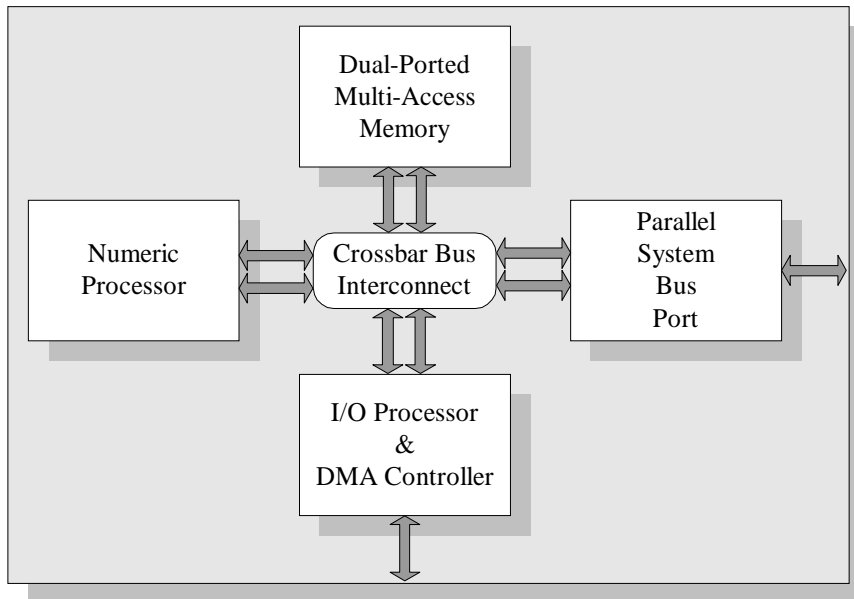
Harvard Architecture
(dual memory)



Super Harvard Architecture
(dual memory, cache,
I/O Processor)

Processor: SHARC

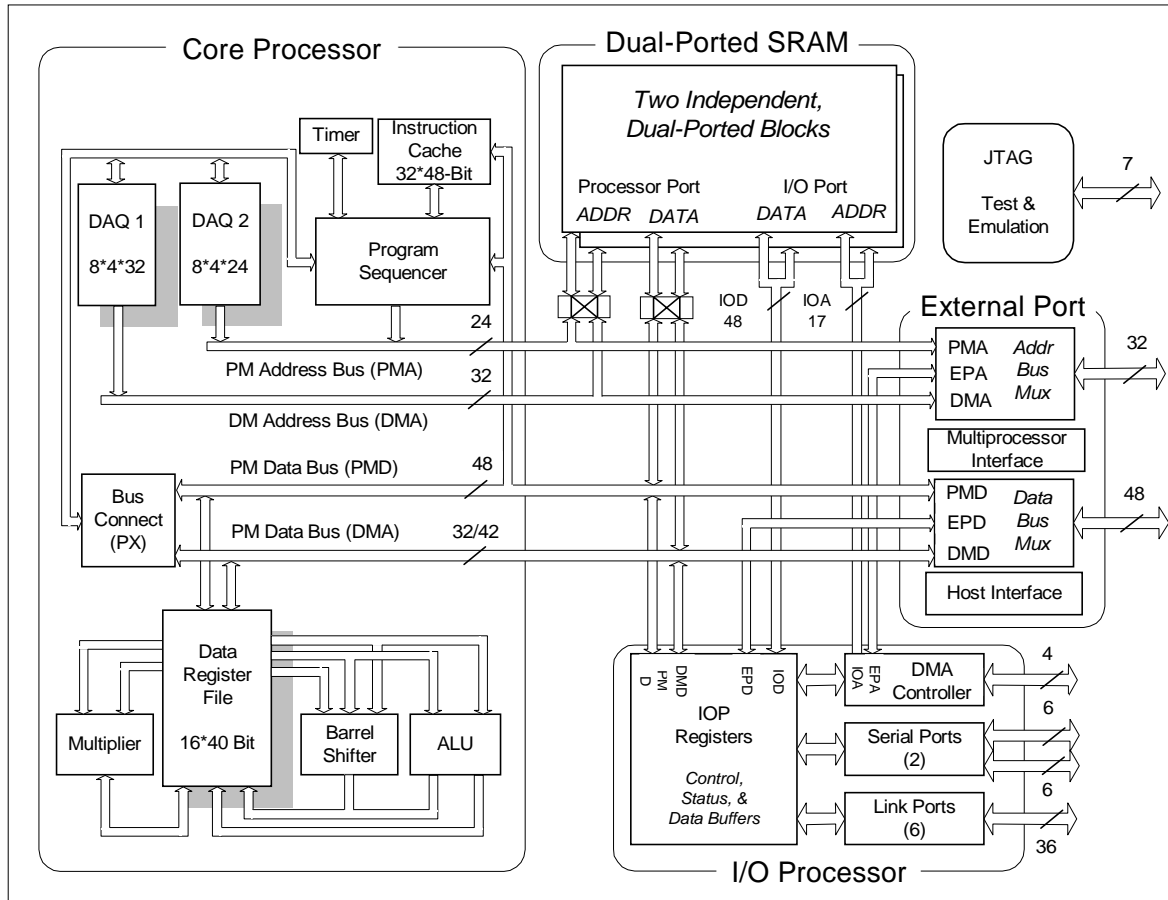
Embedded DSP: SHARC Architecture



- A **crossbar switch** connecting the core numeric processor to an independent I/O processor, dual-ported memory, and parallel system bus port.

SHARC:
Super Harvard Architecture

Embedded DSP: SHARC Architecture



Main Technical data:

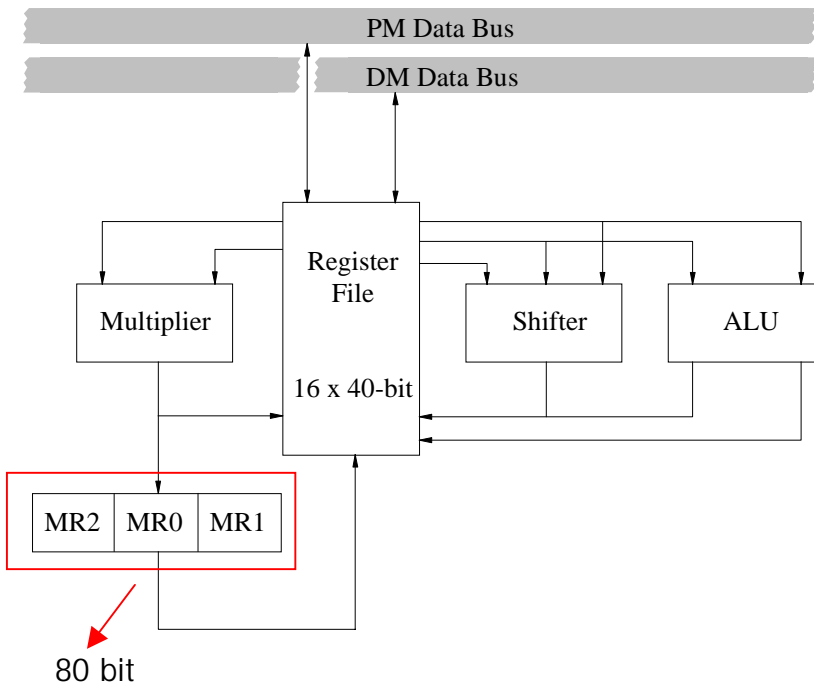
- 32-Bit IEEE FP unit (Multiplie, ALU, and Shifter)
- Data Register File
- Data Address Generator (DAG1, DAG2)
- Program Sequencer with Instruction Cache
- Interval Timer
- Dual-Ported SRAM
- External Port for Interfacing to Off-Chip Mermory & Peripherals
- Host-Port & Multiprocessor Interface
- Serial Ports
- Link Ports
- JTAG Test Access Port

Embedded DSP: Why Floating Point ?

- A digital signal processors data format detemines the ability to handle signals of various
 - precisions
 - dynamic range
 - signal-to-noise ratios
- Dynamic Range:
 - Compression and Decompression algorithms have operated on known bandwidth.
 - Adaptive filtering and imaging are two applications rquiring wide dynamic range.
- Signal-to Noise Ratio:
 - Radar and sonar or speech recognition require wide dynamic range in order to decrease signals from noisy environment.
- In general, 32-bit floating point DSPs are easier to use and allow a quicker and cheaper time-to-market (no scaling !)
- Consistency with IEEE of workstations is a real benefit.

Embedded DSP: Core Processor

- The **core processor** consists of
 - Three computation units **ALU**, **MAC** with a fixed point accumulator, **SHIFTER**
 - three formats: 32-bit fixed point, **32-bit floating-point (IEEE)**, 40-bit floating-point
 - ALU performs a standard set of arithmetic and logic operations in fixed and floating point formats.
 - Multiplier performs floating-point and fixed-point multiplications as well as fixed-point multiply/add and multiply/subtract operations.
 - The shifter performs logical and arithmetic shifts, bit manipulation, field deposit and extraction and exponent derivation operations on 32-bit operands.
 - The computation units perform single-cycle operations (no computation pipeline). The output of any unit may be the input of any unit on the next cycle. In a multifunction computation, the ALU and multiplier perform independent, simultaneous operations.
 - A **register file** is used to transfer data between the computation units and the data buses and for temporary storage of intermediate result. Two sets (**foreground/background**) with each 16 register with 32-bit.



Embedded DSP: Computation Units

- Numeric processing arranged in parallel
- General example fixed-point/floating-point:
 - $F0=F1 * F2$ for floating-point multiply, $R0=R1 * R2$ for fixed-point multiply

ALU Instructions (43)

$R_n = R_x + R_y$	$F_n = F_x + F_y$
$R_n = R_x - R_y$	$F_n = F_x - F_y$
$R_n = (R_x + R_y) / 2$	$F_n = (F_x + F_y) / 2$
COMP(R_x, R_y)	COMP(F_x, F_y)
$R_n = R_x + 1$	$F_n = \text{ABS}(F_x + F_y)$
$R_n = R_x - 1$	$F_n = \text{ABS}(F_x - F_y)$
$R_n = -R_x$	$F_n = -F_x$
$R_n = \text{ABS } R_x$	$F_n = \text{ABS } F_x$
$R_n = \text{PASS } R_x$	$F_n = \text{PASS } F_x$
$R_n = R_x \text{ AND } R_y$	$F_n = \text{RND } F_x$
$R_n = R_x \text{ OR } R_y$	$F_n = \text{LOGB } F_x$
$R_n = R_x \text{ XOR } R_y$	$F_n = \text{ABS } F_x$
$R_n = \text{NOT } R_x$	$F_n = \text{MIN}(F_x, F_y)$
$R_n = \text{MIN}(R_x, R_y)$	$F_n = \text{MAX}(F_x, F_y)$
$R_n = \text{MAX}(R_x, R_y)$	$F_n = \text{RSQRTS } F_x$
.....

MAC Instructions (26)

$R_n = R_x * R_y$ (..)	$F_n = F_x * F_y$
$\text{MRF} = R_x * R_y$ (..)	
$\text{MRB} = R_x * R_y$ (..)	
$R_n = \text{MRF} + R_x * R_y$ (..)	
$R_n = \text{MRB} + R_x * R_y$ (..)	
$\text{MRF} = \text{MRF} + R_x * R_y$ (..)	
$\text{MRB} = \text{MRF} + R_x * R_y$ (..)	
$\text{MRF} = 0$	
$\text{MRB} = 0$	
$\text{MRF} = R_n$	
$\text{MRB} = R_n$	
... =RND ...(..)	
... =SAT ...(..)	
.....	

(..):
 S = Signed input
 U = Unsigned input
 I = Integer input
 F = Fractional input
 FR = Fractional inputs,
 rounded output
 SF = Default format for 1
 -input operation
 SSF = Default format for 2
 -input operations

Embedded DSP: Computation Units

- Shifter operates on 32-bit fixed-point operands.
 - Shift and rotates from off-scale left to off-scale-right
 - Bit manipulation including bit set, clear, toggle and test
 - Bit field manipulation including extract and deposit
 - support for fixed point to floating point conversion

SHIFTER Instructions (36)

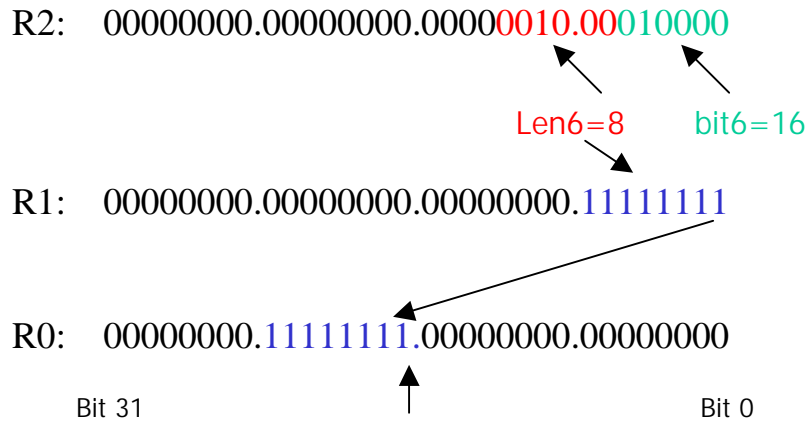
Rn=LSHIFT Rx BY Ry	Rn=BSET Rx BY Ry
Rn=LSHIFT Rx BY <data>	Rn=BSET Rx BY <data>
Rn=Rn OR LSHIFT Rx BY Ry	Rn=BTGL Rx BY Ry
Rn=Rn or LSHIFT Rx BY <data>	Rn=BTGL Rx BY <data>
Rn=ASHIFT Rx BY Ry	Rn=BCLR Rx BY Ry
Rn=ASHIFT Rx BY <data>	Rn=BCLR Rx BY <data>
Rn=Rn OR ASHIFT Rx BY Ry	BTST Rx BY Ry
Rn=Rn or ASHIFT Rx BY <data>	BTST Rx BY <data>
Rn=ROT Rx BY Ry	Rn=FDEP Rx BY Ry
Rn=ROT Rx BY <data>	Rn=FDEP Rx BY <data>
Rn=BCLR Rx BY Ry	Rn=FEXT Rx BY Ry
Rn=BCLR Rx BY <data>	Rn=FEXT Rx BY <data>

Embedded DSP: Computation Units

Field deposit instruction: R0=FDEP R1 BY R2

R1 = 0x000000FF

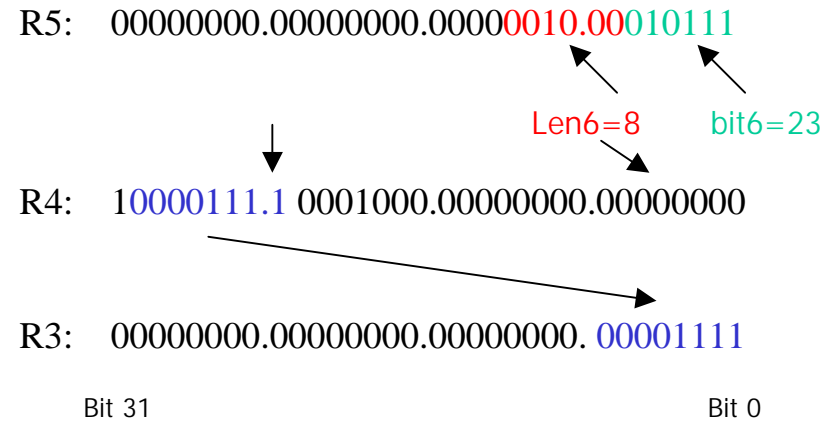
R2 = 0x00000210



Field extract instruction: R3=FEXT R4 BY R5

R4 = 0x87880000

R5 = 0x00000217



Embedded DSP: Computation Units

Multifunction Instructions with MAC/ALU:

$R_a = R_x + R_y$, $R_s = R_x - R_y$;

$R_m = R(3-0) * R(7-4)$ (SSFR),
 $MRF = MRF + R(3-0) * R(7-4)$ (SSFR),
 $R_m = MRF + R(3-0) * R(7-4)$ (SSFR),
 $MRF = MRF - R(3-0) * R(7-4)$ (SSFR),
 $R_m = MRF - R(3-0) * R(7-4)$ (SSFR),

$R_a = R(11-8) + R(15-12)$;
 $R_a = R(11-8) - R(15-12)$;
 $R_a = (R(11-8) - R(15-12))/2$;

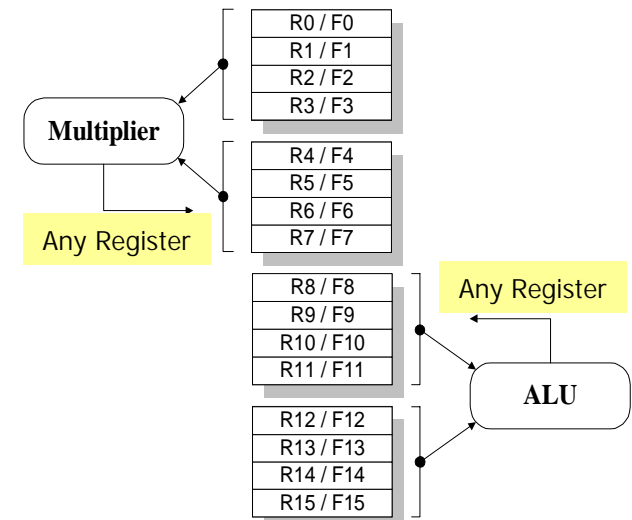
Floating-Point Multiplication and ALU Operation

$F_m = F(3-0) * F(7-4)$,

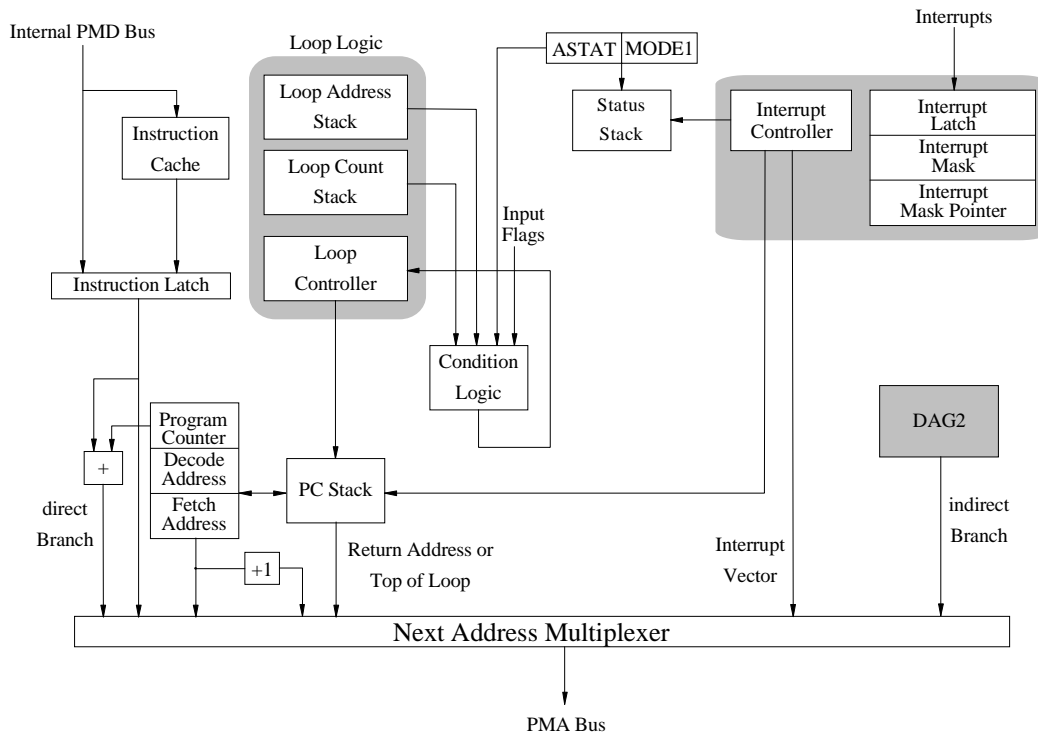
$F_a = F(11-8) + F(15-12)$;
 $F_a = F(11-8) - F(15-12)$;
 $F_a = \text{FLOAT } R(11-8) \text{ BY } R(15-12)$;
 $F_a = \text{ABS } (F(11-8), F(15-12))$;
 $F_a = \text{MAX } F(11-8) + F(15-12)$;
 $F_a = \text{MIN } F(11-8) + F(15-12)$;

Multiplication and Dual Add/Subtract

$R_m = R(3-0) * R(7-4)$ (SSFR), $R_a = R(11-8) + R(15-12)$, $R_s = R(11-8) - R(15-12)$;
 $FM = F(3-0) * F(7-4)$, $F_a = F(11-8) + F(15-12)$, $R_s = F(11-8) - F(15-12)$;



Embedded DSP: Core Processor

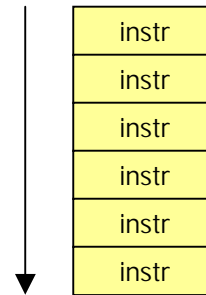


Remember to the pipelined operation:
fetch/decode/execute

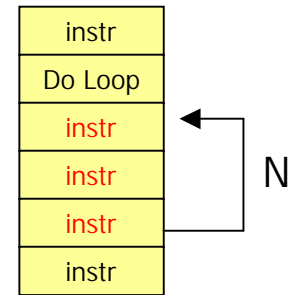
- The **core processor** consists of
 - Program sequencer
 - Address generators
 - DAGs with pre- and postmodify and length register (circular buffer)
 - Instruction cache memory (32*32 bit)
 - Loop Logic with zero overhead
 - Branch Logic
 - Control & Status register
 - Program counter
 - PC stack unit
 - Interrupt controller
 - external: 3, internal: 26 (vectorized by fixed addresses and priority)
 - Programmable timer (timer for periodically interrupts, watch-dog function)

Embedded DSP: Program Sequencer

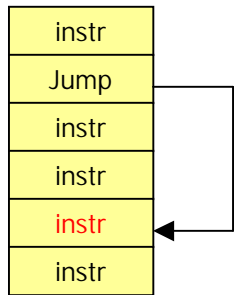
- Program Flow by:
 - Loops, subroutines, jumps, interrupts, idle
 - incrementing the fetch address,
 - maintaining stacks,
 - evaluating conditions,
 - decrementing loop counter
 - calculating new addresses
 - maintaining an instruction cache
 - handling interrupts



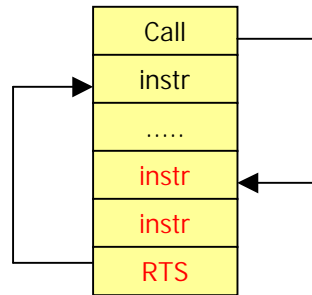
Linear Flow



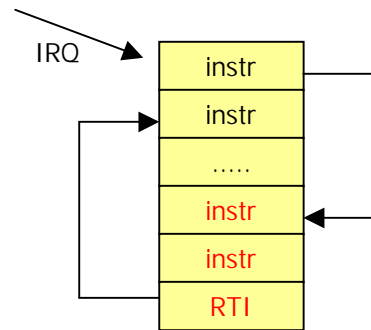
Loop



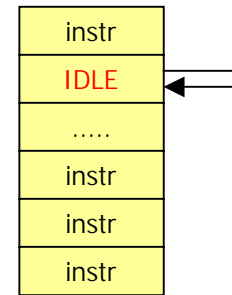
Jump



Subroutine



Interrupt



Subroutine

Embedded DSP: Overview SHARC Instructions

Program Flow Control Instructions

Nr.				Addressing Mode	
8a	If condition	JUMP	<addr24> (PC,<reladdr24>)	(DB) ; (LA) (CI) (DB,LA) (DB,CI)	Direct addressing Relative addressing
8b	If condition	CALL	<addr24> (PC,<reladdr24>)	(DB) ;	Direct addressing Relative addressing
9a	If condition	JUMP	(Md,lc) (PC,<reladdr6>)	(DB) , compute ; (LA) ELSE compute (CI) (DB,LA) (DB,CI)	Indirect ,pre-modify' Relative addressing
9b	If condition	CALL	(Md,lc) (PC,<reladdr6>)	(DB) , compute ; ELSE compute ;	Indirect ,pre-modify' Relative addressing
10	If condition	JUMP	(Md,lc), (PC,<reladdr6>),	, ELSE , compute, DM(la,Mb)=dreg ; compute, dreg=DM(la,Mb)	Indirect ,pre-modify' Relative addressing
11a	If condition	RTS	(DB) , (LR) , (DB,LR) ,	, compute ; ELSE compute	Direct addressing Relative addressing
11b	If condition	RTI	(DB),	, compute ; ELSE compute	Direct addressing
12	LCNTRL=	<data16> ureg	, DO <addr24 (PC,<reladdr24>)	UNTIL LCE ;	Direct addressing Relative addressing
13	DO	<addr24> (PC,<reladdr24>)	UNTIL termination ;		Direct addressing Relative addressing

- Condition and Loop Termination Codes:

		<u>True if</u>
• EQ	ALU equal zero	AZ=1
• LT	ALU less than zero	..
• LE	ALU less than or equal zero	..
• AC	ALU carry	AC=1
• AV	ALU overflow	AV=1
• MV	Multiplier overflow	MV=1
• MS	Multiplier sign	MS=1
• SV	Shifter overflow	SV=1
• SZ	Shifter zero	SZ=1
• ...		
• Complements of the above conditions:		
• NEQ		AZ=0
• GE		..
• GT		..
• NOT AC		AC=0
•		

Examples:

CALL init (DB);

JUMP (M8,I12), R6=R6-1;

IF EQ CALL(PC,17) (DB), ELSE R6=R6-1;

IF NOT GT RTS(DB);

IF SZ RTS, ELSE R0=LSHIFT R1 BY R15;

LCNTR=100, DO fmax UNTIL CE;

LCNTR=R12, DO (PC,16) UNTIL CE;

DO end UNTIL FLAG1_IN

DO (PC,7) UNTIL AC

Embedded DSP: Program Sequencer

Branches (Loops and Calls) can be delayed (DB) or nondelayed:

- If **nondelayed**, the **two instructions** after the branch, which are in the fetch and decode stages, **are not executed**.
- For a call the decode address (the address of the instruction after the call) is the return address.
- During the two no-operations cycles, the first instruction at the branch address is fetched and decoded.

- If **delayed**, the processor **continues to execute two more instructions** while the instruction at the branch address is fetched and decoded.
- In the case of a **call**, the return address is the **third** address after the branch instruction.
- A delayed branch is more efficient, but it makes the code harder to understand because of the instructions between the branch instructions and the actual branch.

- Restrictions: Instructions in the two lines following a delayed branch may not be:
 - Other Jumps, Calls or Returns
 - Pushes or Pops of the PC stack
 - Writes to the PC stack or PC stack pointer
 - DO UNTIL instruction
 - IDLE instruction

Embedded DSP: Program Sequencer

non delayed branch:

```
JUMP label;
(NOP; )
(NOP; )
```

delayed branch:

```
JUMP label (DB)
R1=DM( I0 ,M0 );
R2=PM( I8 ,M8 );
```

Programm :

⋮
Command A
Command B
Jump
Command C
Command D
⋮
Command E
Command F
Command G

Example: Pipeline execution

delayed or not delayed branch ?

Pipeline :

fetch	Jump	C	D	E	F	G
decode	B	Jump	C	empty	E	F
execute	A	B	Jump	empty	empty	E

```
LCNTR=30, DO label1 UNTIL CE;
    F0=DM(I0,M0), F2=PM(I8,M8);
    F1=F0*F2;
label1:  F4=F1+F4;
```

Embedded DSP: Data Addressing

Overview

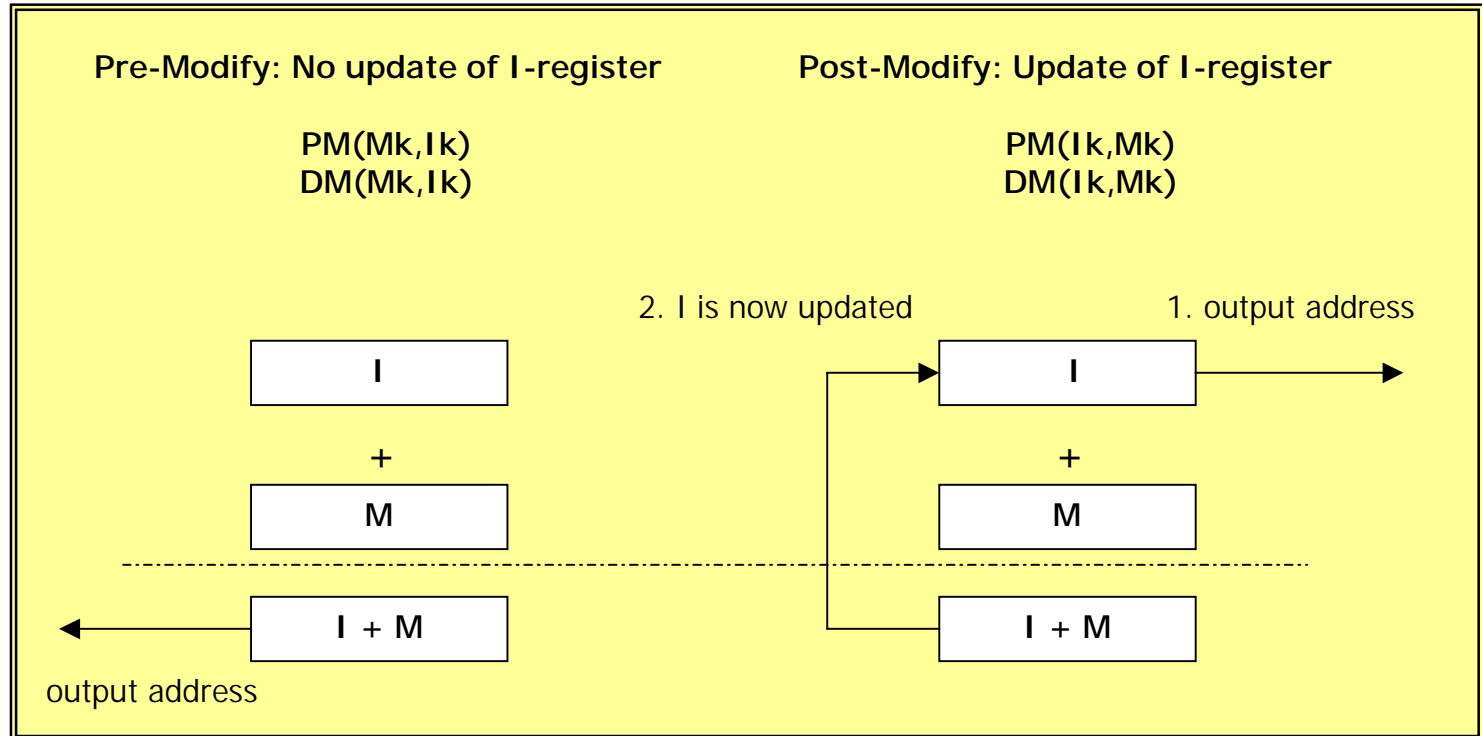
- 2 data address generators (DAG1 and DAG2)
- indirect address access, address indirect by the content of a DAG
- DAG1 addresses 32-bit on data bus DM
- DAG2 addresses 24-bit on program bus PM
- Additional alternate (secondary) register for fast context switching
- Special support for signal processing applications
 - Circular data buffers
 - Bit-reversing
- Each has 4 types of special registers:

• Index (I)	= pointer to memory	[DAG1: I0-I7] [DAG2: I8-I15]
• Modify (M)	= increment/decrement value	[M0-M7] [M8-M15]
• Base (B)	= base address of a circular buffer	[B0-B7] [B8-B15]
• Length (L)	= length of a circular buffer	[L0-L7] [L8-L15]

DAG operation

- address output (pre-modify or post-modify)
- modulo addressing (circular buffers)
- addressing in bit-reverse order

Embedded DSP: Data Addressing



Pre-Modify versus Post-Modify

Modifier Instructions

- **R6 = PM(I12,M11);** Indirect addressing PM-memory with post-modify
 - R6 = content of I12
 - I12 = I12 + M11
- **R6 = PM(M11,I12);** Indirect addressing PM-memory with pre-modify
 - R6 = content of (I12 + M11)
 - I12 is not changed
- **DM(M1,I2) = TCOUNT;** Indirect addressing DM-memory
 - Store TCOUNT in DM address (I2+M1)
- **R2 = DM(0x40000012, I1);** Immediate 32-bit modify: address = I1 + 0x40000012
- **F6 = F1 + F3, PM(I8,0x0A)=ASTAT;** Immediate 6-bit modify: address=I8, I8=I8+0x0A

Instruction Set Reference

- **Compute and Move or Modify** Instructions, which specify a compute operation in parallel with one or two data moves or an index register
- **Program Flow Control** instructions, which specify various types of branches, calls, returns and loops. Some of these instructions may also specify a compute operation and/or a data move
- **Immediate Data Move** instructions, which use immediate instruction fields as operands, or use immediate instruction fields for addressing.
- **Miscellaneous** instructions, such as bit modify and test, no operation and idle

Embedded DSP: Overview SHARC Instructions

Compute & Move or Modify Instructions

Nr.	Instruction	Addressing mode
1	compute , DM(Ia,Mb)=dreg1 , PM(Ic,Md)=dreg2 ; dreg1=DM(Ia,Mb) , dreg2=PM(Ic,Md)	Indirect addressing ,post-modify'
2	If condition compute ;	
3a	If condition compute , DM(Ia,Mb)=ureg ; PM(Ic,Md)=ureg	Indirect addressing ,post-modify'
3b	If condition compute , DM(Ma,Ib)=ureg ; PM(Mc,Ic)=ureg	Indirect addressing ,pre-modify'
3c	If condition compute , ureg=DM(Ia,Mb) ; ureg=PM(Ic,Md)	Indirect addressing ,post-modify'
3d	If condition compute , ureg=DM(Ma,Ib) ; ureg=PM(Mc,Ic)	Indirect addressing ,pre-modify'
4a	If condition compute , DM(<data6>)=dreg ; PM(<data6>)=dreg	Immediate addressing mode
4b	If condition compute , DM(<data6>,Ia)=dreg ; PM(<data6>,Ic)=dreg	Immediate addressing mode ,pre-modify'
4c	If condition compute , dreg=DM(<Ia,<data6>) ; dreg=PM(Ic,<data6>)	Immediate addressing mode ,pre-modify'
4d	If condition compute , dreg=DM(<data6>,Ia) ; dreg=PM(<data6>,Ic)	Addressing mode: ,Immediate-pre-modifier'
5	If condition compute , ureg1=ureg2 ;	Register transfer
6a	If condition shiftimm , DM(Ia,Mb)=dreg ; PM(Ic,Md)=dreg	Indirect addressing ,post-modify'
6b	If condition shiftimm , dreg=DM(Ia,Mb) ; dreg=PM(Ic,Md)	Indirect addressing ,post-modify'
7	If condition shiftimm , Modify(Ia,Mb) ; Modify(Ic,Md)	Modify address (Ia=Ia+Mb) (Ic=Ic+Md)

Embedded DSP: Overview SHARC Instructions

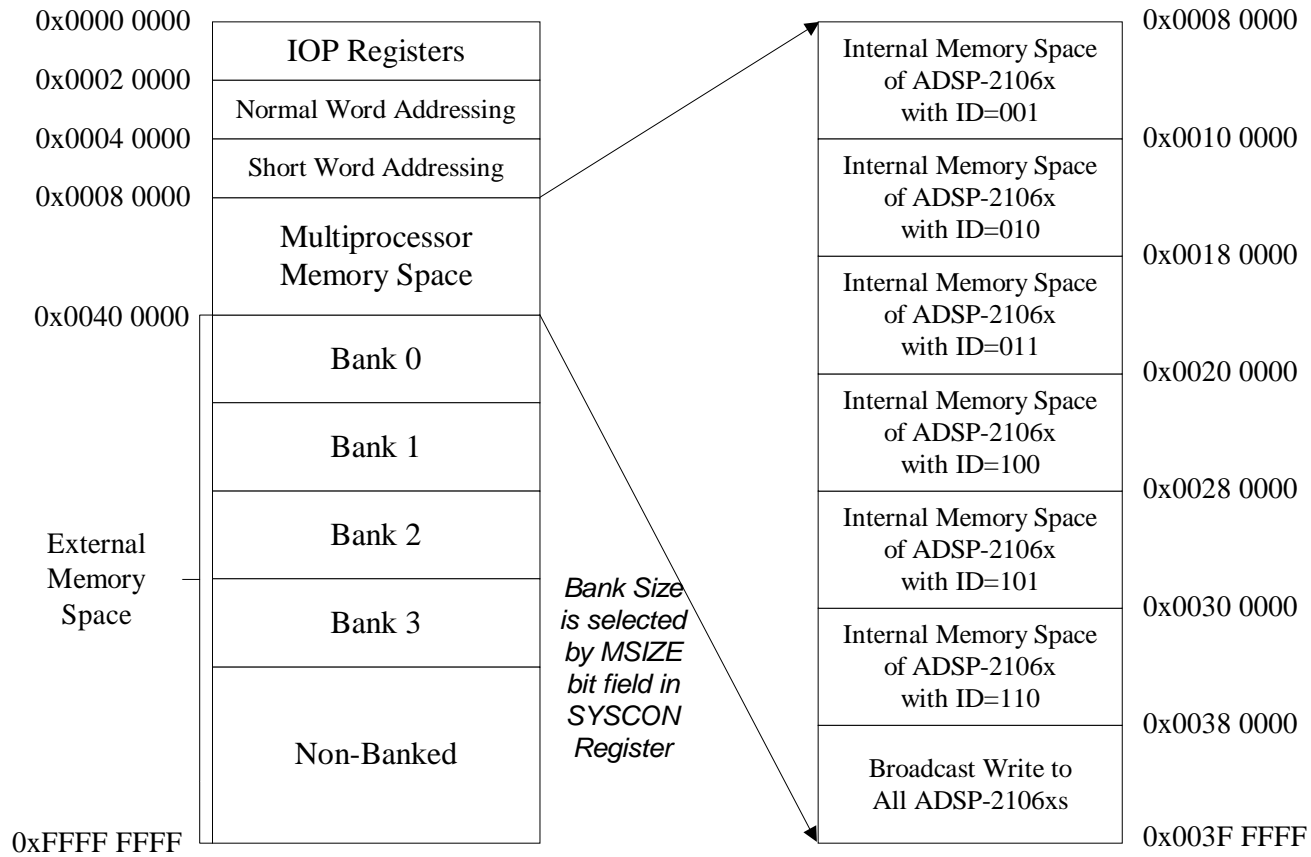
Immediate Move Instructions

Nr.		Addressing Mode
14a	DM(<addr32>) = ureg ; PM(<addr24>)	Direct addressing
14b	ureg= DM(<addr32>) ; PM(<addr24>)	Direct addressing
15a	DM(<data32>) = ureg ; PM(<data24>)	Direct addressing
15b	ureg= DM(<data32>) ; PM(<data24>)	Direct addressing
16	DM(la,Mb) = <data32> ; PM(lc,Md)	Direct addressing
17	ureg = <data32> ;	Direct addressing

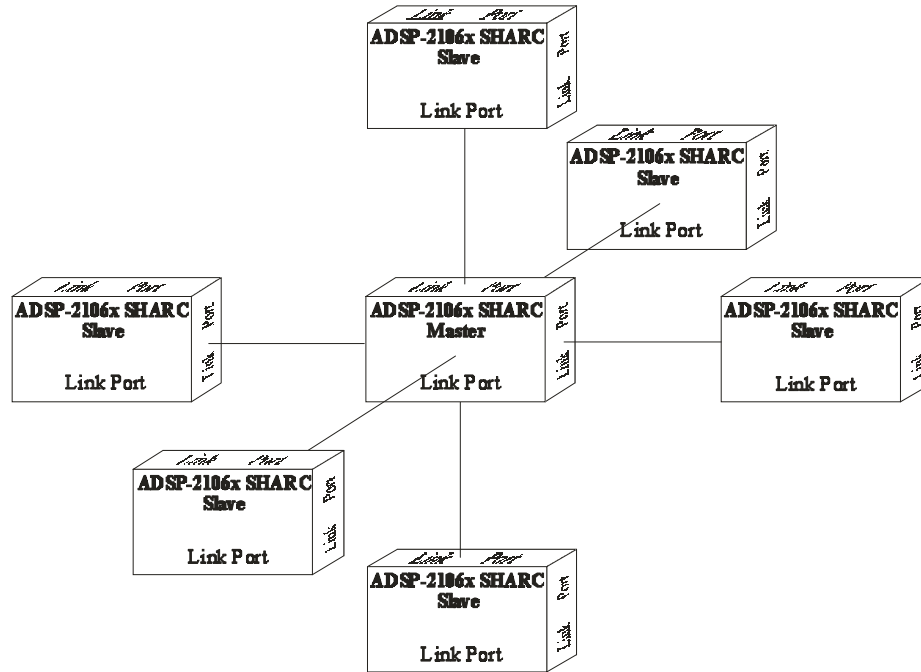
Miscellaneous Instructions

Nr.	Instruction
18	BIT SET sreg <data32> ; CLEAR TGL TST XOR
19a	MODIFY (la,<data32> ; (lc,<data24>
19b	BITREV (la,<data32> ; (lc,<data24>
20	PUSH LOOP , PUSH STS , PUSH PCSTK , FLUSH ; POP POP POP CACHE
21	NOP ;
22	IDLE ;
23	CJUMP function (DB) ; (PC,<reladdr24>
24	RFRAME ;

Embedded DSP: Internal Memory

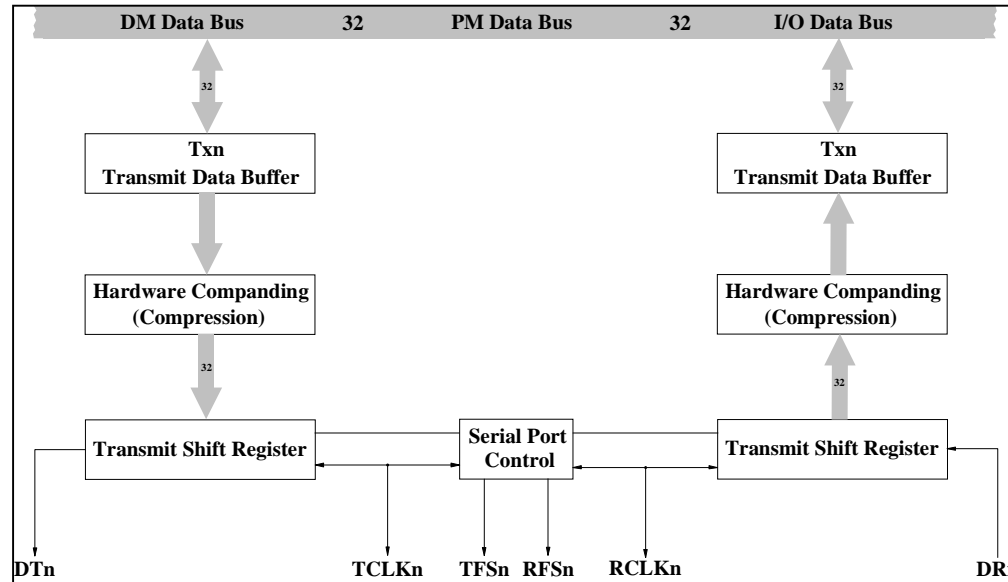


Embedded DSP: Multiprocessing by LINK ports

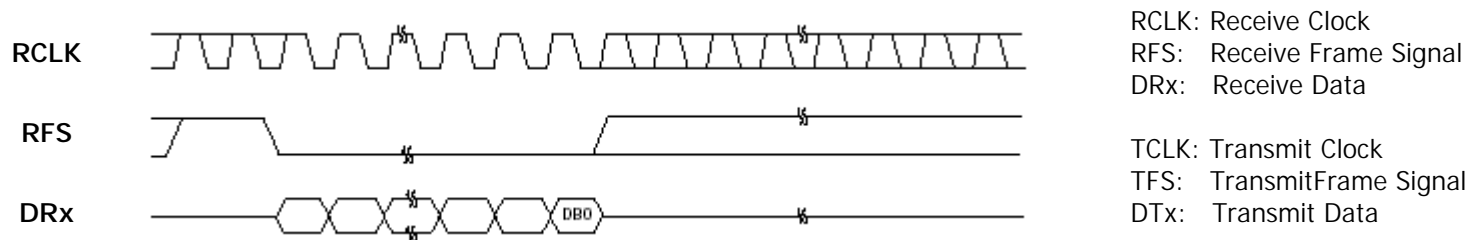


- Link ports for high speed point-to-point data transfers to other processors
- SHARC: 21062: 6 Link ports
- 4 -bit ports
- transmit and receive
- double buffer register
- Handshake signals REQ & ACK
- Single cycle or DMA
- Interrupts
- Fully asynchron to the core
- only 6 pins !

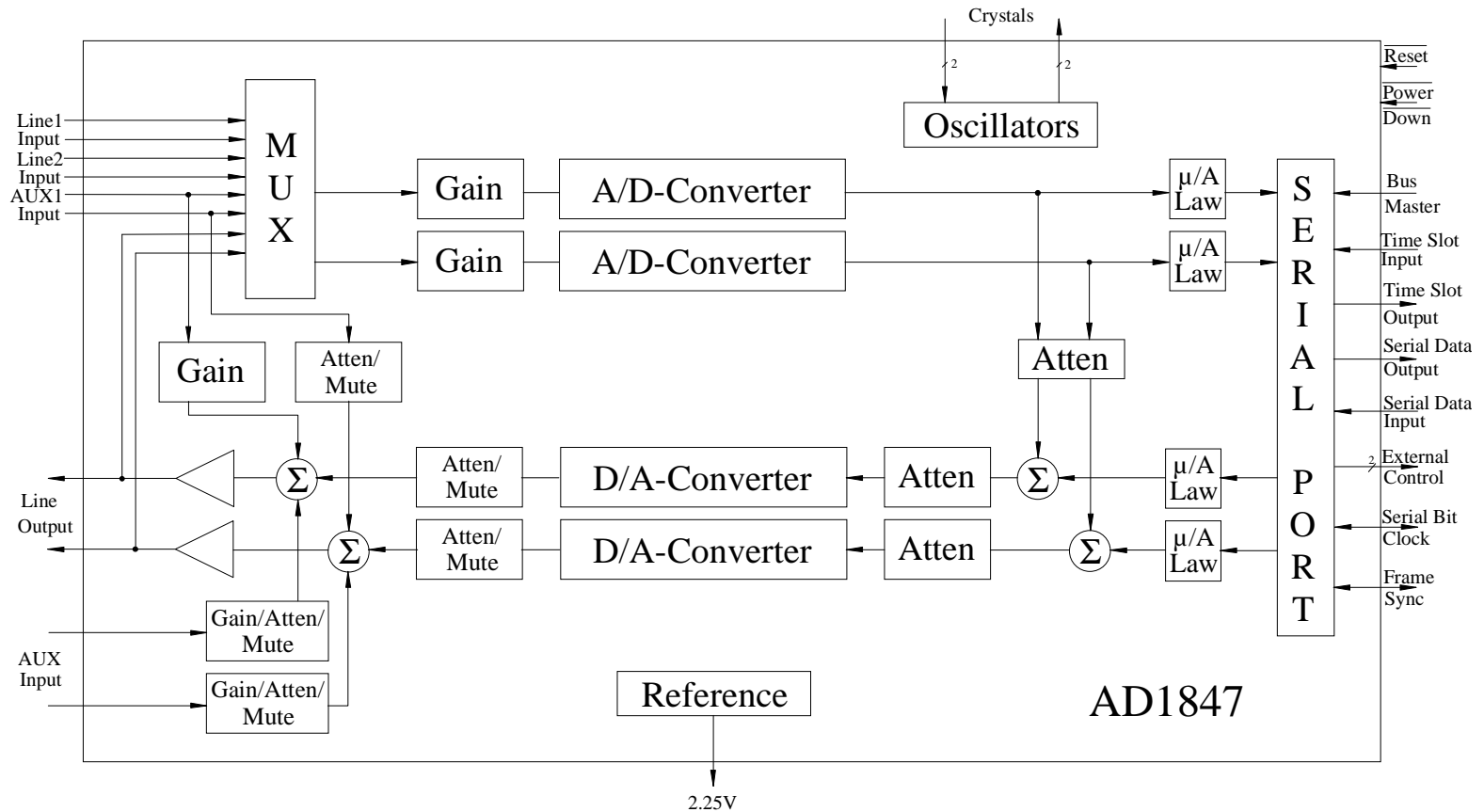
Embedded DSP: Serial Ports



3 Wire Interface



Embedded DSP: Serial Port to a CODEC



Time and Frequency Considerations of SSI based Converters

Serial Interface			Processor		Analog Signal
Cycle time [us]	Transfer rate [M bit/sec]	Transfer time 16 Bit [us]	Cycle time [ns]	Count of Instructions	Frequency
1	1	16	0.025	640	< 32 kHz
0.1	10	1.6	0.025	64	< 312 kHz
0.05	20	0.8	0.025	32	< 625 kHz
0.03	33	0.48	0.025	19	< 1 MHz

Result:

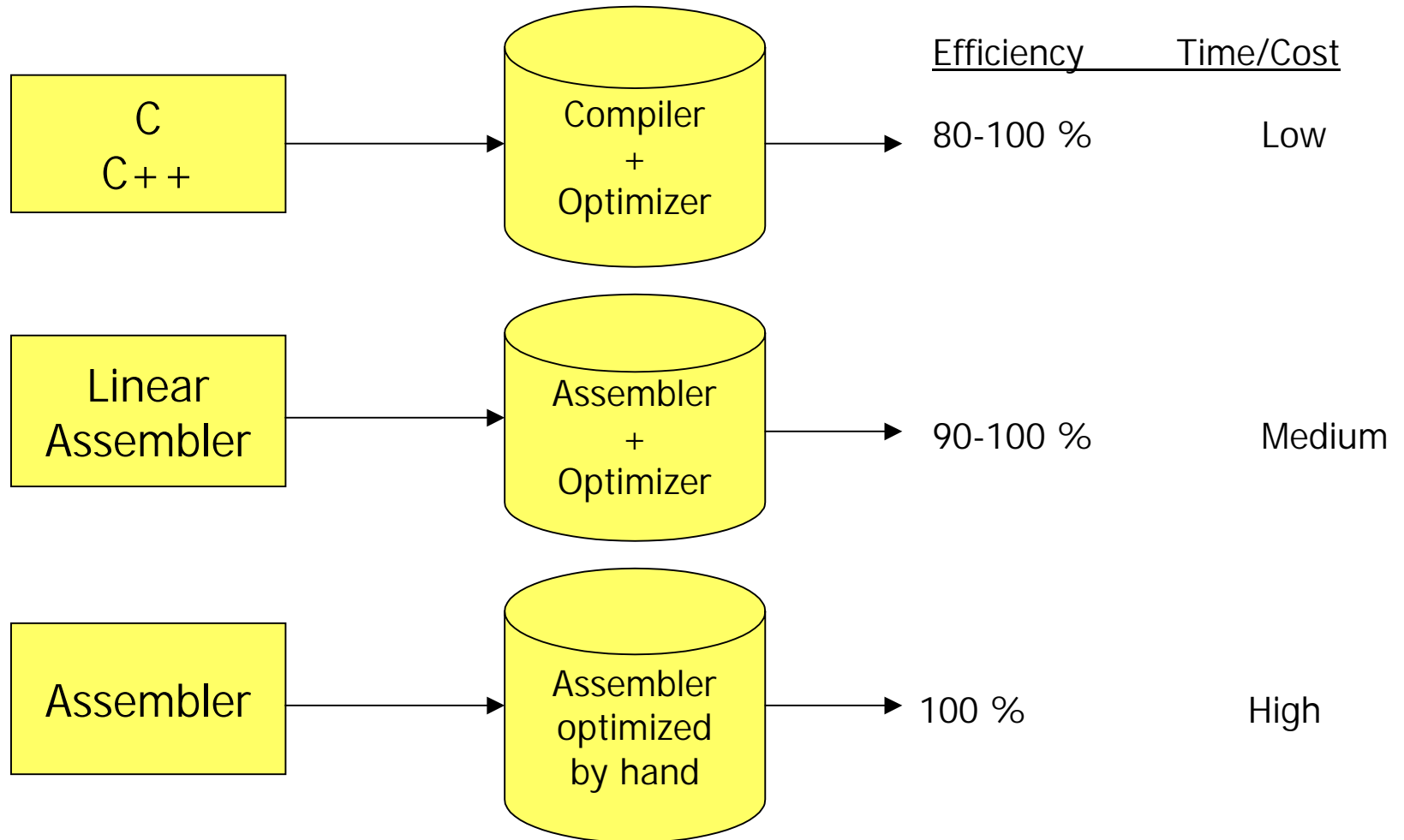
- Use the highest possible data rate on SSI interface
- Use DMA transfer to relieve the processor
- Use interrupt technology for event triggering
- Typical situation for audio processing
 - 48 kHz sample rate
 - 20.8 us / sample
 - 20.8 us = interrupt
 - 40 MHz DSP <-> at least 832 single instructions

Embedded DSP: Software

Outline

- Programming Methods
 - Text editor based
 - Visual DSP Studio
- Projects
- Working with C
- Working with Assembler

Embedded DSP: Software

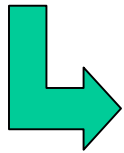


```

void CCpyFifoToMem()
{
  for (i = 0; i < TsCnt; i++) {
    value = FIFO(4) >> 16;
    rbuf[i] = abs(value);
  }
}

```

C-Program



ASM

18

```

.extern      _abs;
.global     _CCpyFifoToMem;
_CCpyFifoToMem:
  modify(i7,-1); nop;
  r2=0;
  dm(_i)=r2;

_L$249:
  r2=dm(_i);
  r4=dm(_TsCnt);
  comp(r2,r4);
  if ge jump (pc, _L$250) ;
  r2=dm(4456452);
  r4=ashift r2 by -16;
  dm(-2,i6)=r4;
  r2=dm(_i);
  r4=_rbuf;
  r8=r4+r2;
  i4=r8;
  r2=dm(-2,i6);
  r4=abs r2;
  dm(i4,m5)=r4;

_L$251:
  r2=dm(_i);
  r4=r2+1;
  r2=r4;
  dm(_i)=r2;
  jump (pc, _L$249) ;

_L$250:
_L$248:
  i12=dm(-1,i6);
  jump (m14,i12) (DB);
  nop;
  RFRAME;

```

Assembler hand coded

```

.SEGMENT /pm  seg_pmco;
.global
_ASMFastCpyExtAbsFifoToMem;
.extern      _rbuf;
.extern      _TsCnt;
_ASMFastCpyExtAbsFifoToMem:
  entry;
  puts=r0;
  puts=r2;
  r0=i1; puts=r0;
  r0=i4; puts=r0;

  i4=0x440004;
  i1=_rbuf;
  r0=dm(_TsCnt);
  lcntr=r0, DO fifo_abs UNTIL LCE;
  r2=dm(i4,m5);
  r2=abs r2;
  r2=ashift r2 by 0xfffff0;
  fifo_abs: dm(i1,m6)=r2;

  r0=gets(1); i4=r0;
  r0=gets(2); i1=r0;
  r2=gets(3);
  r0=gets(4);
  alter(4);
  exit;
.endseg;

```

4

Assembler hand coded

```
.SEGMENT /pm seg_pmco;
.global
_ASMFastCpyExtAbsFifoToMem;
.extern _rbuf;
.extern _TsCnt;
_ASMFastCpyExtAbsFifoToMem:
entry;
puts=r0;
puts=r2;
r0=i1; puts=r0;
r0=i4; puts=r0;

i4=0x440004;
i1=_rbuf;
r0=dm(_TsCnt); /* TsCnt */
lcntr=r0, DO fifo_abs UNTIL LCE;
    r2=dm(i4,m5);
    r2=abs r2;
    r2=ashift r2 by 0xfffff0;
fifo_abs: dm(i1,m6)=r2;

r0=gets(1); i4=r0;
r0=gets(2); i1=r0;
r2=gets(3);
r0=gets(4);
alter(4);
exit;
.endseg;
```

4



Assembler hand coded (++)

```
.SEGMENT /pm seg_pmco;
.global _ASMFastAbs;
.extern _rbuf;
.extern _TsCnt;
_ASMFastAbs:
entry;
puts=r0; puts=r2; puts r4;
r0=i1; puts=r0;
r0=i4; puts=r0;

i4=0x440004;
i1=_rbuf;
r0=dm(_TsCnt); /* TsCnt = TsCnt/2 */
r2=dm(i4,m5);

lcntr=r0, DO abs2 UNTIL LCE;
    r2 = ashift r2 by 0xfffff0; r4=dm(i4,m5);
    r2 = abs r2;
    r4 = ashift r4 by 0xfffff0; dm(i1,m6)=r2;
    r4 = abs r4; r2=dm(i4,m5);
abs2: dm(i1,m6)=r4;

r0=gets(1); i4=r0;
r0=gets(2); i1=r0;
r4=gets(3);
r2=gets(4);
r0=gets(5);
alter(5);
exit;
.endseg;
```

5

```

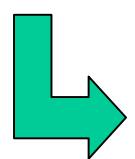
void CCpyFifoToMem()
{
    int value, *ip0;

    ip0=rbuf;

    for (i=0; i<TsCnt; i++)
    {
        value = mk_SRAM2(4) >> 16;
        *ip0++ = abs(value);
    }
}

```

C-Program



ASM

20

```

•extern      _abs;
_CCpyFifoToMem:
    modify(i7,-3);
    r2=i0;
    dm(-4,i6)=r2;
    r2=_rbuf;
    dm(-3,i6)=r2;
    r2=0;
    dm(_i)=r2;
_L$249:
    r2=dm(_i);
    r4=dm(_TsCnt);
    comp(r2,r4);
    if ge jump (pc, _L$250) ;
    r2=dm(4456452);
    r4=ashift r2 by -16;
    dm(-2,i6)=r4;
    i4=dm(-3,i6);
    i0=i4;
    modify(i0,m6);
    i2=i0;
    dm(-3,i6)=r2;
    r2=dm(-2,i6);
    r4=abs r2;
    dm(i4,m5)=r4;

```

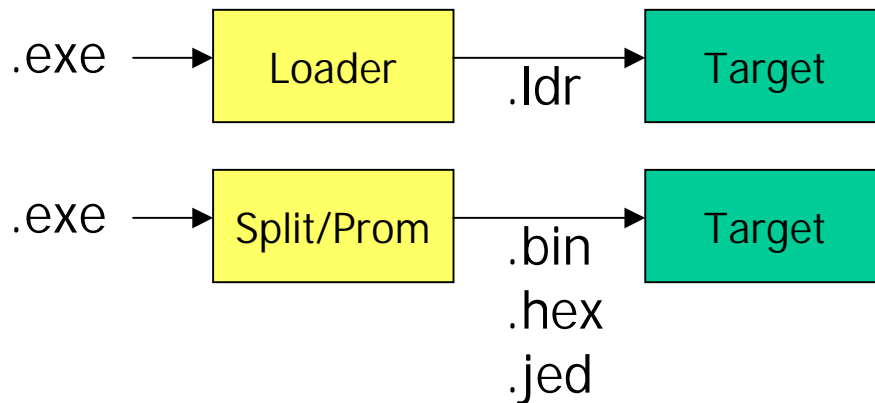
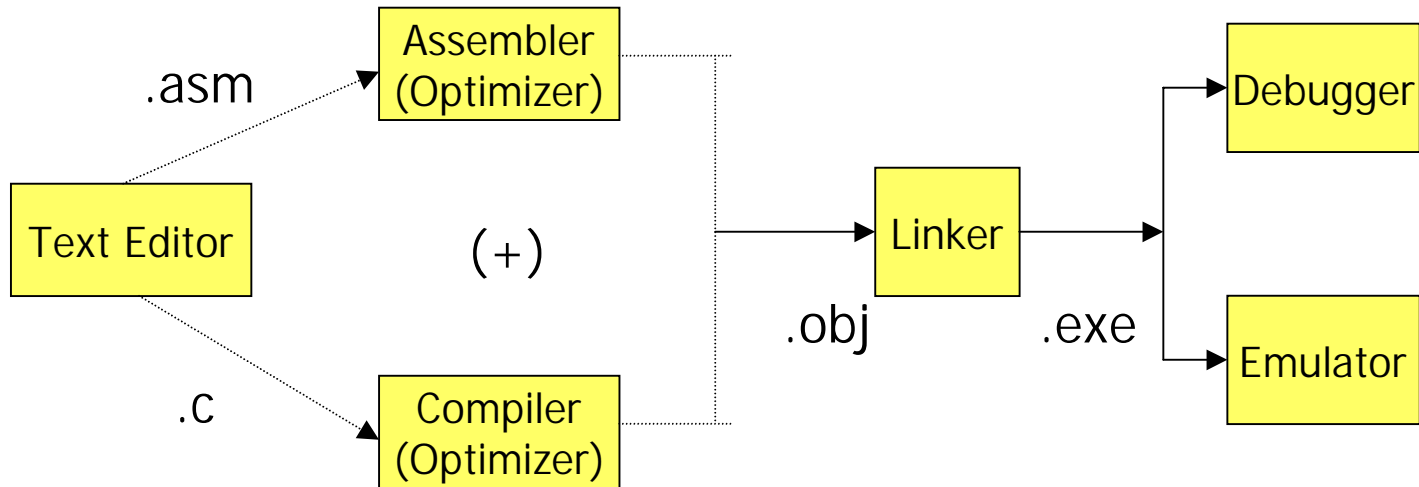
```

_L$251:
    r2=dm(_i);
    r4=r2+1;
    r2=r4;
    dm(_i)=r2;
    jump (pc, _L$249) ;
_L$250:
_L$248:
    i12=dm(-1,i6);
    i0=dm(-4,i6);
    jump (m14,i12) (DB);
    nop;
RFRAME;

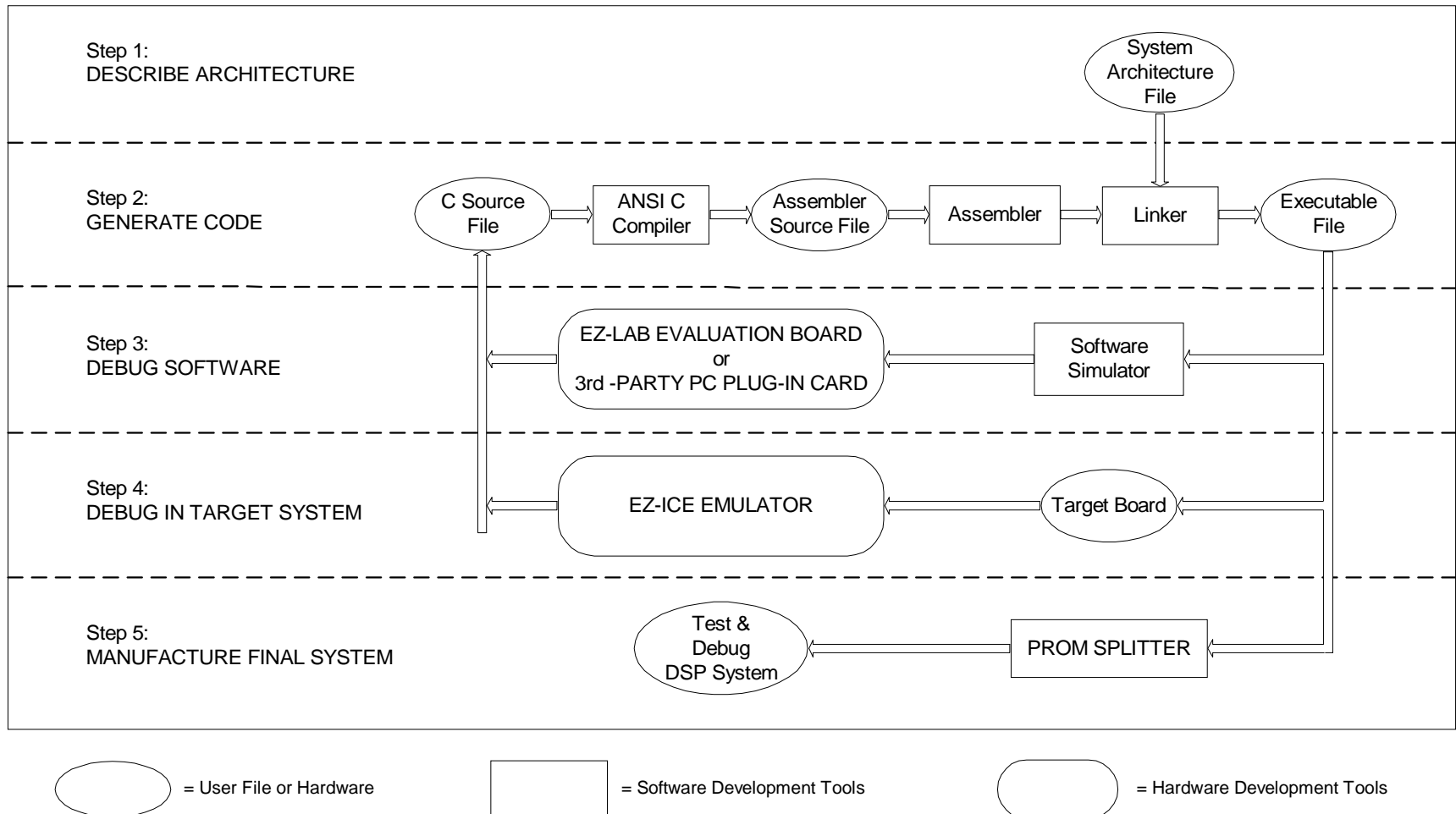
```

Embedded DSP: SHARC Software Tools

Method 1:



Development Environment with different modules during the 90th



SIMULATOR

The screenshot displays the ADSP-2106x Simulator interface with the following components:

- Memory (asm, pc):** Shows assembly code for memory addresses from 0002037f to 0002038f. The instruction at 0002038f, `dm(i7,m7)=pc;`, is highlighted.
- CBUG (sh_pc_dc.c):** Shows the source code for `init_components()` starting at line 80. The code includes `enable_interrupts()`, a `while(1)` loop, and a `switch` statement for `board_interrupt`. The console below shows a restart, a hit breakpoint at line 80, a user halt, and another hit breakpoint at line 80.
- Memory : 2 (hex, off):** Shows memory values for addresses 0002448b to 00024491, including `_frequenz: 000001c3`, `_lower: 00000002`, `_upper: 00000023`, and various `_erg` registers.
- Active Register File:** Lists registers R0 through R15 with their current values in hexadecimal.
- Status Bar:** Displays "Target Halted".

Embedded DSP: Method2: SHARC-Visual DSP

Since 1999:

- Windows based software development environment for Analog Devices Digital Signal Processors:
 - An integrated development environment with support for editing programs, managing projects, and controlling build tools.
 - A source level, object oriented debugger with support for DSP simulation and emulation.
 - Context-sensitive help for the Windows-based development environment.
 - Online access to all documentation for the products via -pdf files.
 - Boot-Loader and generation of HEX and S files for EPROMs

Embedded DSP: Method2: SHARC-Visual DSP

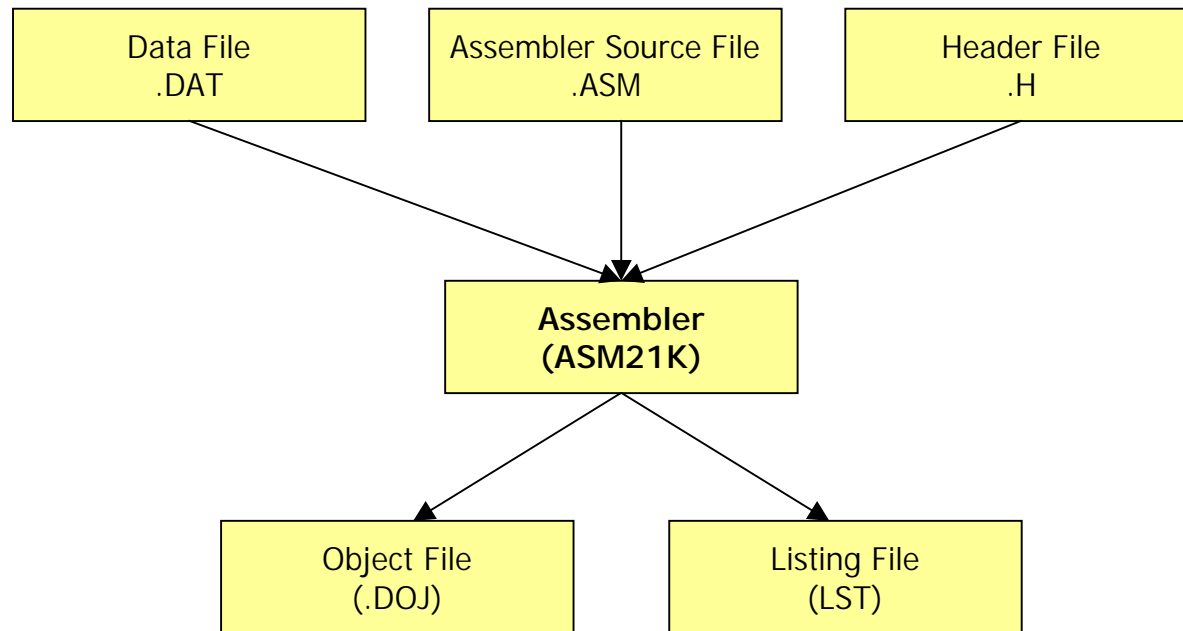
- VisualDSP supports many file formats:
 - Source files
 - Assembly Source Files (*.ASM)
 - Assembly Initialization Data Files (*.DAT)
 - Preprocessor Header Files (*.H)
 - Linker Description File (*.LDF)
 - ASCII text files that contain commands for the linker scripting language, which holds target information
 - Linker Command File (*.TXT)
 - Build (Processed) Files
 - Assembler Object Files (*.DOJ)
 - Archiver File (*.DLB)
 - Linker Executable File (*.DXE, *.SM, *.OVL, *.DLO)
 - Linker Memory Map (*.MAP)
 - Loader HEX-Files, ASCII-Files, Binary-Files (*.LDR)
 - Splitter Motorola S-Record Files (*.S_#)
 - Splitter Hex Format Files (*.H_#)
 - Splitter Byte-Stacked Format Files (.STK)

Embedded DSP: Method2: SHARC-Visual DSP

- Debugger Files
 - Provide input to the debugger to define support for simulation or emulation of the program.
 - The debugger supports all the executable files types produced by the linker (.DXE, .SM, .OVL, .DLO).
 - To simulate I/O, the debugger supports the data file formats (.DAT) from the assembler, the loadable file formats from the loader (.LDR), and the PROM formats from the splitter (.S_#, .H_#, .STK)

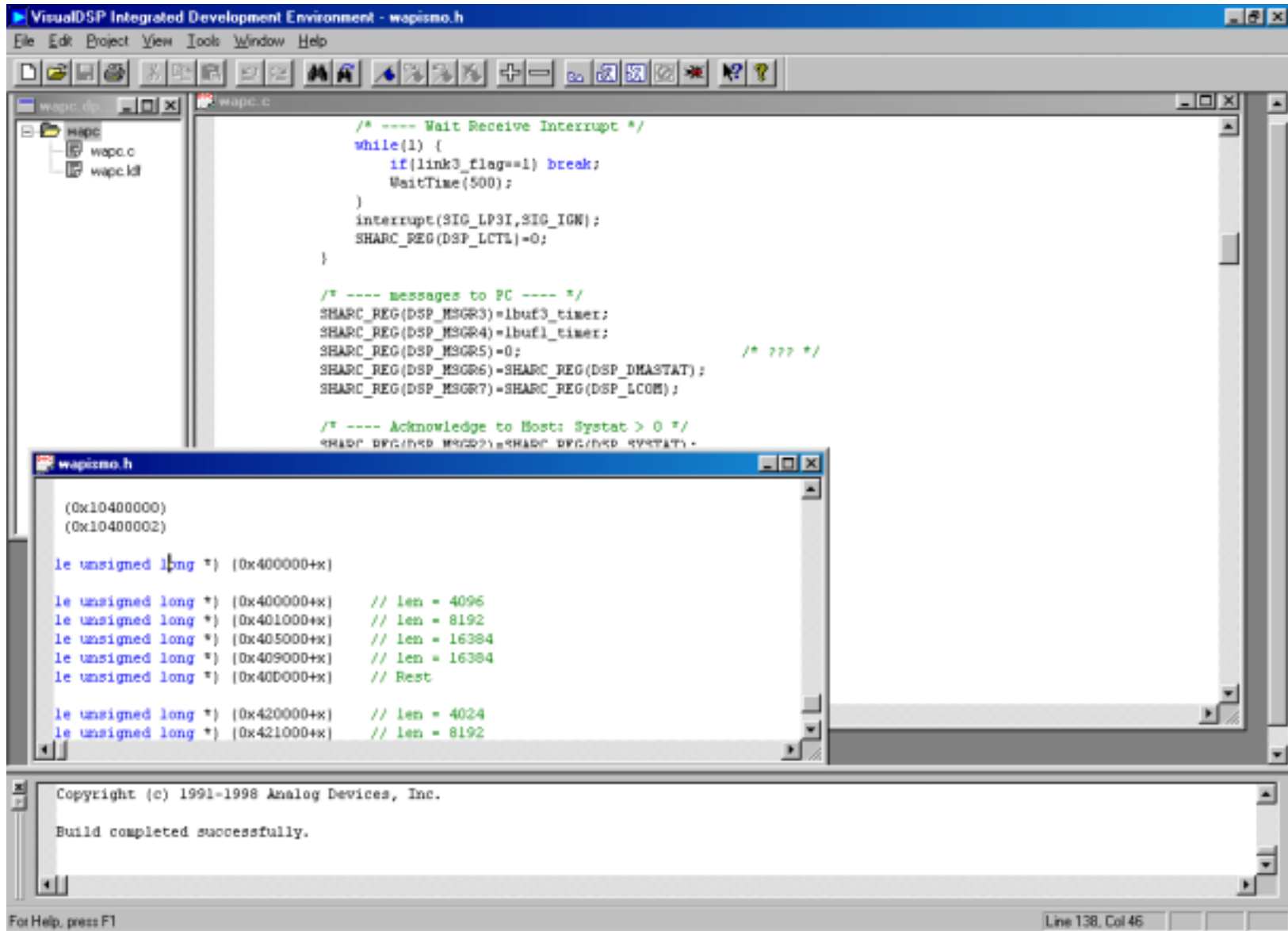
Embedded DSP: SHARC Software Tools

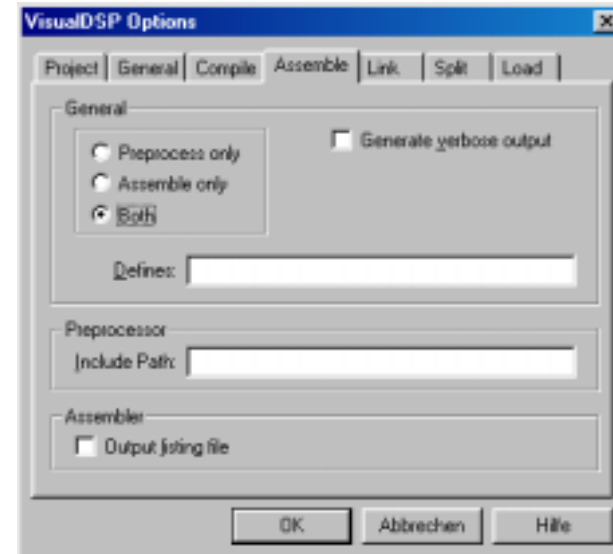
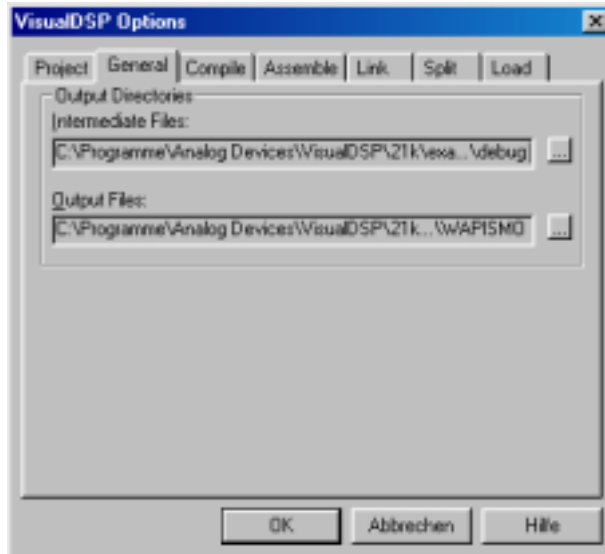
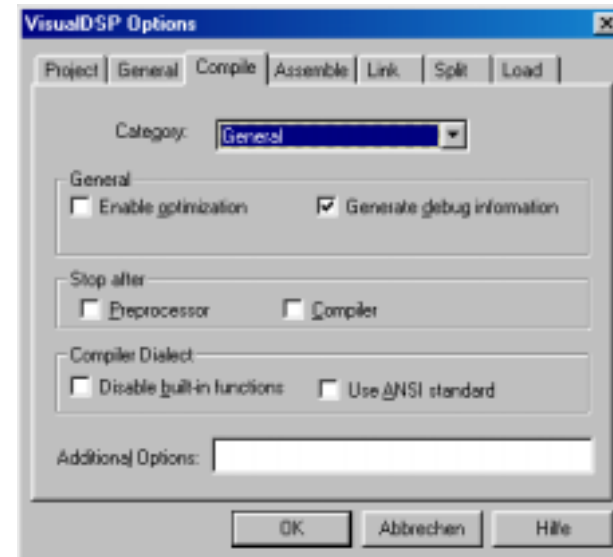
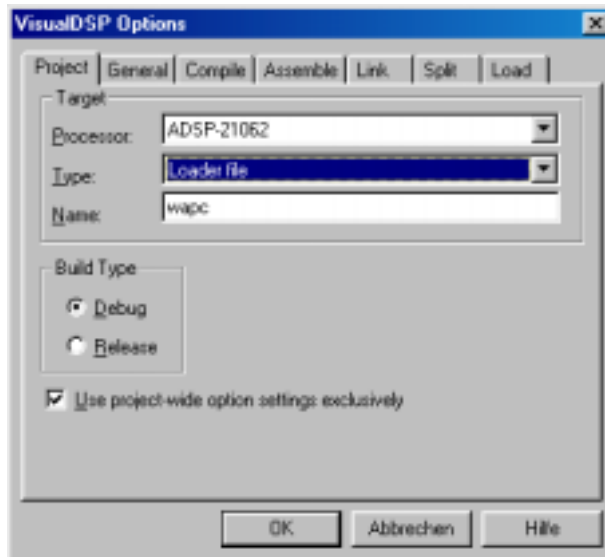
- Assembler Input & Output Files

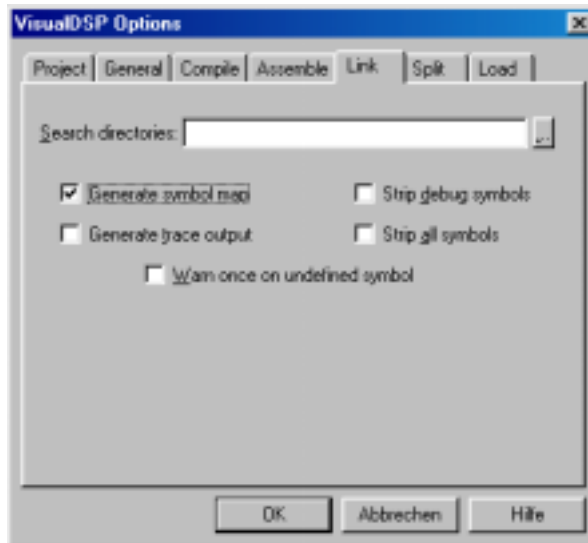
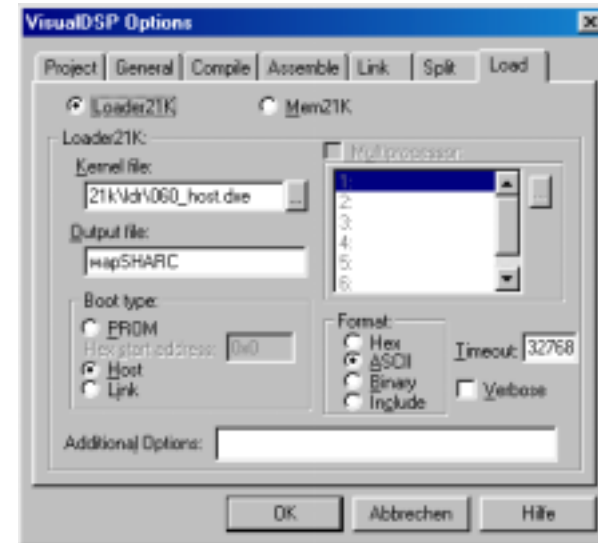
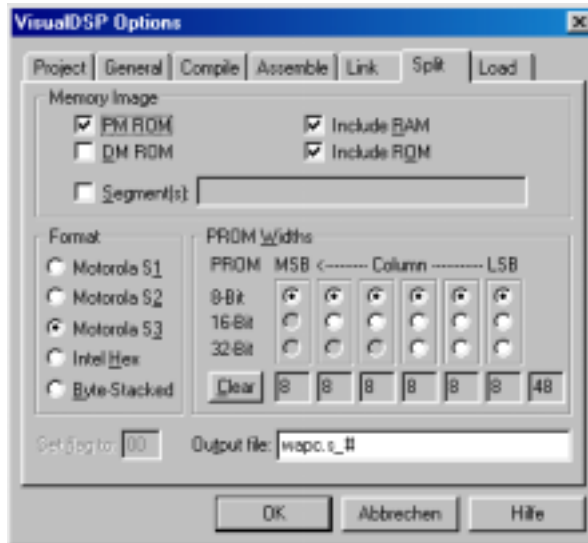


Embedded DSP: SHARC Software Tools

- C-Compiler is realized by an additional Runtime Header Model
 - Preserves structures for register usage, runtime stack model and stack heap.
 - Need for a special memory model in the architecture file
 - Uses the following default segments:
 - seg_rth (Interrupt table/runtime header)
 - seg_pmco (PM code)
 - seg_pmda (PM data)
 - seg_dmda (DM data)
 - seg_heap (heap space)
 - seg_stack (stack space)







Standard Functions

CHARACTER & STRING MANIPULATION

atoi	convert string to integer
bsearch	binary search of array
isalnum	detect alphanumeric character
isalpha	detect alphabetic character
isctrl	detect control character
isdigit	detect decimal digit
isgraph	detect printable character
islower	detect lowercase character
isprint	detect printable character
ispunct	detect punctuation character
isspace	detect whitespace character
isupper	detect uppercase character
isxdigit	detect hexadecimal digit
memchr	find first occurrence of char
memcpy	copy characters
strcat	concatenate strings
strcmp	compare strings
strerror	get error message
strlen	string length
strncmp	compare characters
strrchr	find last occurrence of char
strstr	find string within string
strtok	convert string to tokens
system	sent string to operating system
tolower	change uppercase to lowercase
toupper	change lowercase to uppercase

MATH OPERATIONS

abs	absolute value
acos	arc cosine
asin	arc sine
atan	arc tangent
atan2	arc tangent of quotient
cabsf	complex absolute value
cexpf	complex exponential
cos	cosine
cosh	hyperbolic cosine
cot	cotangent
div	division
exp	exponential
fmod	modulus
log	natural logarithm
log10	base 10 logarithm
matadd	matrix addition
matmul	matrix multiplication
pow	raise to a power
rand	random number generator
sin	sine
sinh	hyperbolic sine
sqrt	square root
srand	random number seed
tan	tangent
tanh	hyperbolic tangent

Standard Functions

PROGRAM CONTROL

abort	abnormal program end
calloc	allocate / initialize memory
free	deallocate memory
idle	processor idle instruction
interrupt	define interrupt handling
poll_flag_in	test input flag
set_flag	sets the processor flags
timer_off	disable processor timer
timer_on	enable processor timer
timer_set	initialize processor timer

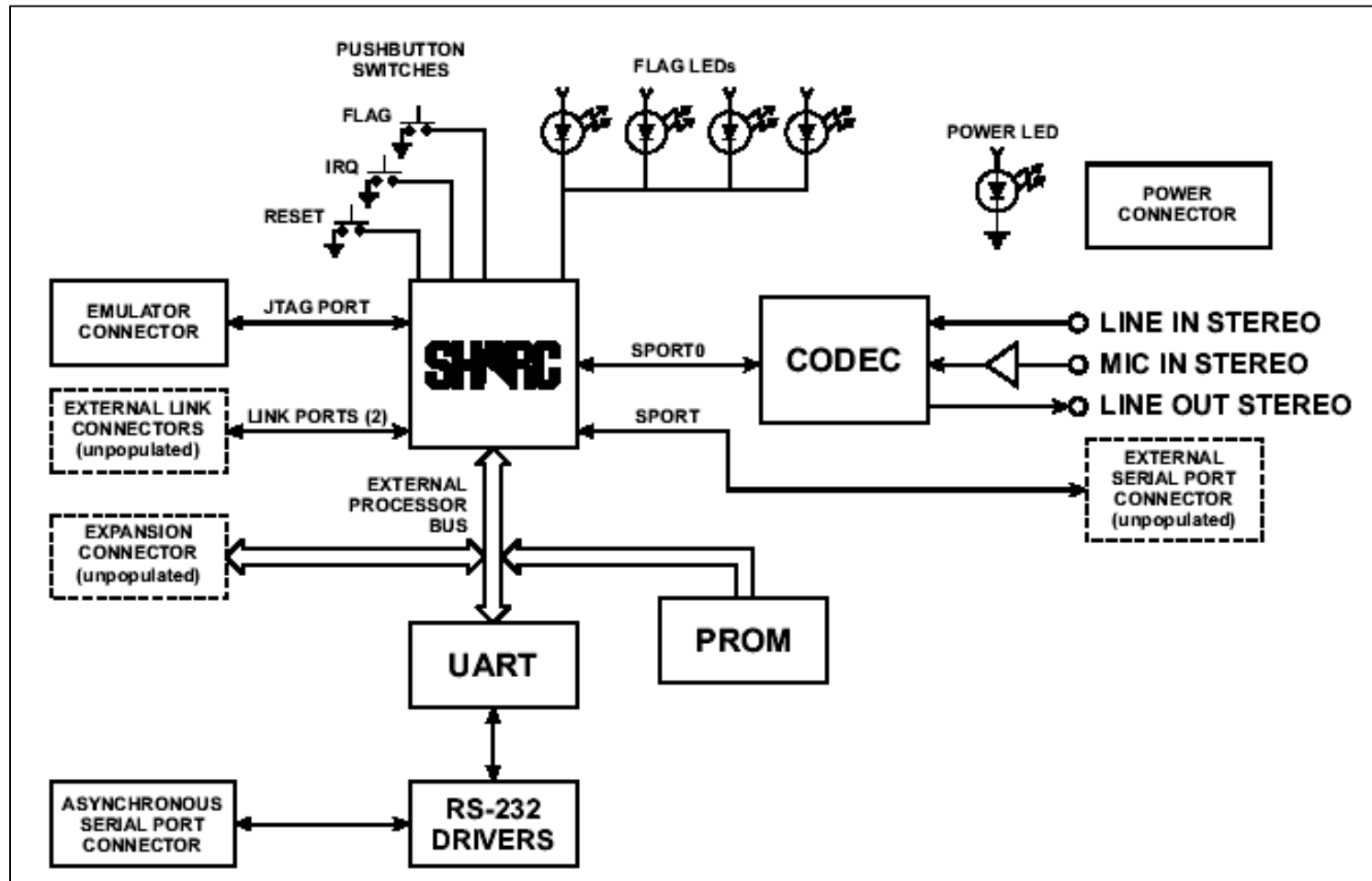
SIGNAL PROCESSING

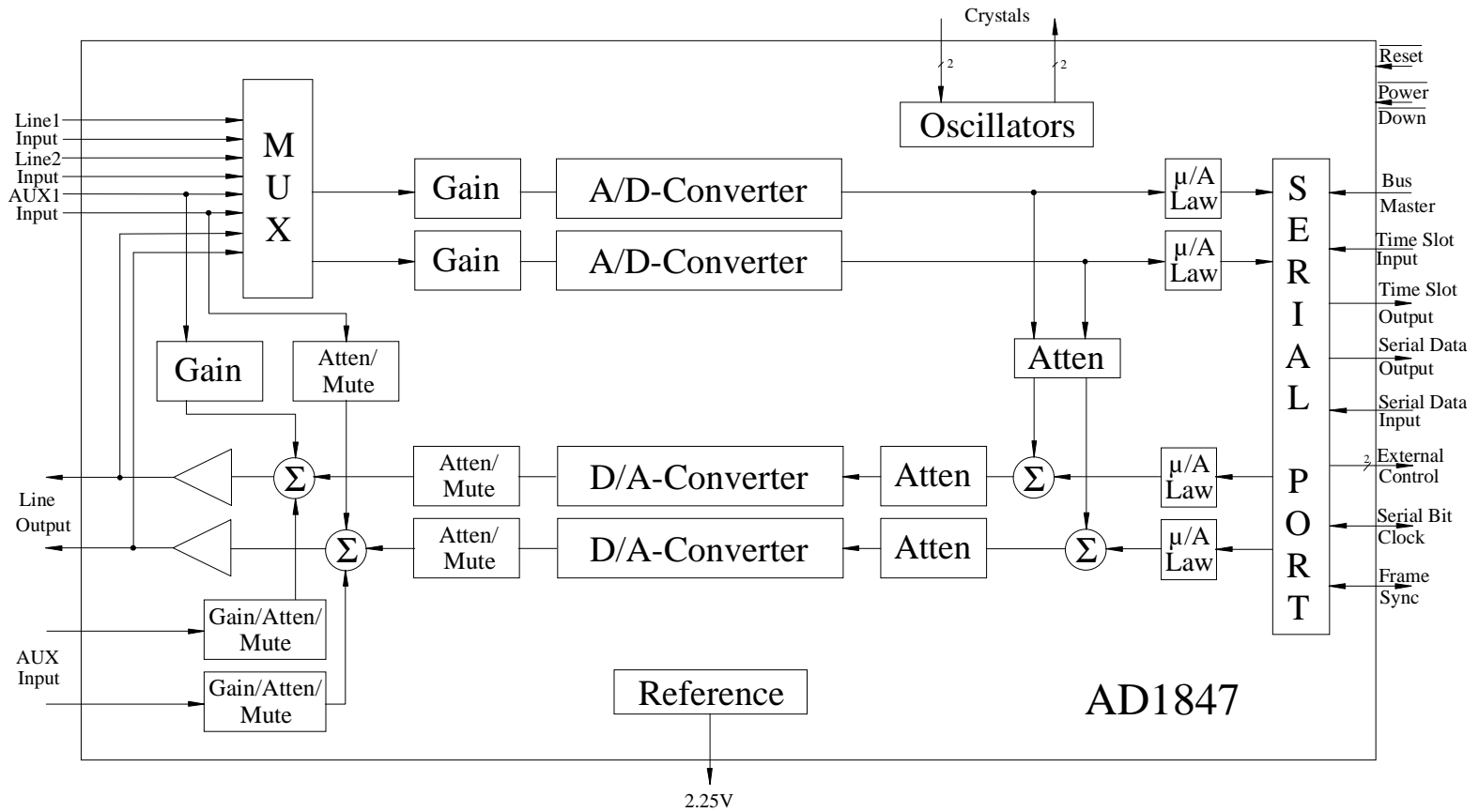
a_compress	A-law compressing
a_expand	A-law expansion
autocorr	autocorrelation
biquad	biquad filter section
cfftN	complex FFT
crosscorr	cross-correlation
fir	FIR filter
histogram	histogram
ifftN	inverse complex FFT
iir	IIR filter
mean	mean of an array
mu_compress	mu law compression
mu_expand	mu law expansion
rfftN	real FFT
rms	rms value of an array

EZ-KIT

PRODUCT	Memory	Notes
AD1460	4 Mbit ×4	Quad-SHARC, Four ADSP-21060's in the same module; provides an incredible 480 MFLOPS in only 2.05"×2.05"×0.16".
ADSP-21160M	4 Mbit	New! Features Single Instruction Multiple Data (SIMD) core architecture; optimized for multiprocessing with link ports, 64 bit external bus, and 14 channels of DMA
ADSP-21060	4 Mbit	Power house of the family; most memory; link ports for high speed data transfer and multi-processing
ADSP-21062	2 Mbit	Same features as the ADSP-21060, but with less internal memory (SRAM), for lower cost
ADSP-21061	1 Mbit	Low cost version used in the EZ-KIT Lite; less memory & no link ports; additional features in DMA for the serial port
ADSP-21065L	544 kbit	A recent addition to the family; fast and very low cost (\$10). Will attract many fixed point applications to the SHARC family
ADSP-21020	-0-	Oldest member of the family. Contains the core processor, but no on-chip memory or I/O interface. Not quite a SHARC DSP.

EZ-KIT





Start of the Embedded DSP Mini Project
next monday at IZFP.