



Estimating the Performance of Cache Replacement Policies

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- 1 Introduction
 - Locality & Caches
 - Problem & Prior Work

- 2 Stochastic Model
 - Contributions
 - Model Construction

- 3 Evaluation
 - Precision & Runtime
 - Related Work

- 4 Conclusions



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- Typical programs do not behave randomly, they exhibit
 - ▶ temporal locality
 - ▶ spatial localityin their memory access behavior
- Caches try to exploit locality
- Have great influence on system performance



Cache Parameters

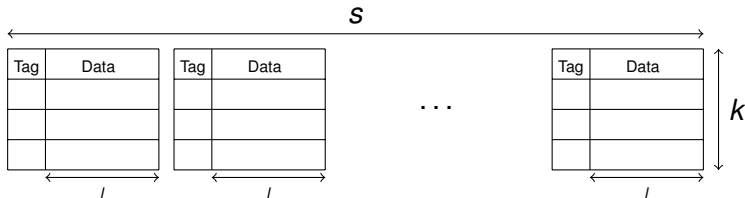
Line size l Size of blocks stored in cache

Associativity k Number of lines one block may be stored in

Number of sets s Each set comprises k lines

Capacity Determined by $s \cdot k \cdot l$

Replacement policy *next slide*





- The replacement policy determines the cache line whose contents are replaced upon a cache miss
- Examples:
 - **LRU** Least recently used
 - **FIFO** First-in first-out
 - **PLRU** Approximation on LRU
 - **MRU** Most recently used

...



- Cache performance depends on (locality of) executed program
- There are lots of design possibilities

Problem

Given a program, determine the best cache configuration.

Stack Histograms (Mattson et al. 1970)



- Concisely capture locality in memory access behavior
- **Age** of cacheable element e :
number of accesses on different elements since last access on e
- **Stack distance** of a memory access:
age of the accessed element just before the access
- **Stack histogram**:
frequency distribution of stack distances

Example

Reference string	a	b	a	c	b	b	c	a
Stack distance	∞	∞	1	∞	2	0	1	2
Stack distance d	0	1	2	3	≥ 4			
Rel. frequency f_d	$\frac{1}{8}$	$\frac{2}{8}$	$\frac{2}{8}$	0	$\frac{3}{8}$			

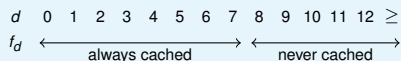


Mattson et al. (1970)

- LRU always retains the k most recently used elements
- Exactly the ages $0, \dots, k - 1$ are cached

Example

Associativity $k = 8$



$$\Rightarrow \text{Miss ratio} = 1 - \sum_{i=0}^7 f_i$$



Problem:

- In general, the ages of cached elements vary during execution
- Number and shape of combinations of ages depend on policy

Idea: determine

- all possible combinations (state space)
- probability of each state
- probability of a cache miss in each state



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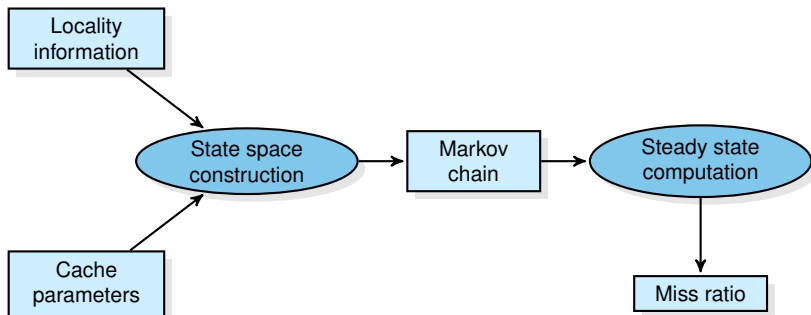
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Our Approach



- Static method
- Given ...
 - ▶ a concise characterization of the program's locality
 - ▶ parameters of the cache architecture
- ... we construct a Markov chain to estimate the miss ratio





- **Replacement policy** is taken into account
- **Policy tables**
 - ▶ uniformly represent replacement policies
 - ▶ ease automatic model construction
- **History stack histograms**
 - ▶ extend well-known stack histograms
 - ▶ increase precision of estimates



- Uniform representation for a class of replacement policies
- Class contains LRU, PLRU, FIFO, MRU
- Eases automatic model construction

π_0	0	1	2	3	4	5	6	7
π_1	0	1	2	3	4	5	6	7
π_2	0	1	2	3	4	5	6	7
π_3	0	1	2	3	4	5	6	7
π_4	0	1	2	3	4	5	6	7
π_5	0	1	2	3	4	5	6	7
π_6	0	1	2	3	4	5	6	7
π_7	0	1	2	3	4	5	6	7
π_m	1	2	3	4	5	6	7	0

8-way FIFO

π_0	1	2	3	4	5	6	7	0
π_1	0	2	3	4	5	6	7	1
π_2	0	1	3	4	5	6	7	2
π_3	0	1	2	4	5	6	7	3
π_4	0	1	2	3	5	6	7	4
π_5	0	1	2	3	4	6	7	5
π_6	0	1	2	3	4	5	7	6
π_7	0	1	2	3	4	5	6	7
π_m	1	2	3	4	5	6	7	0

8-way LRU

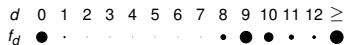
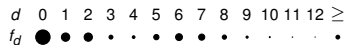
- *Policy state* is virtual order on cache lines
[3, 2, 1, 0, 4, 5, 6, 7]
- *Updates* are specified by permutations
 - ▶ π_i for cache-hit on position i
 - ▶ π_m for cache-miss

$$[3, 2, 1, 0, 4, 5, 6, 7] \xrightarrow{\pi_4} [3, 2, 1, 0, 5, 6, 7, 4]$$

History Stack Histograms



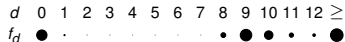
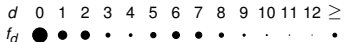
Circle plots of stack histograms



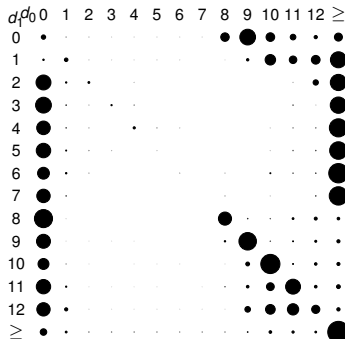
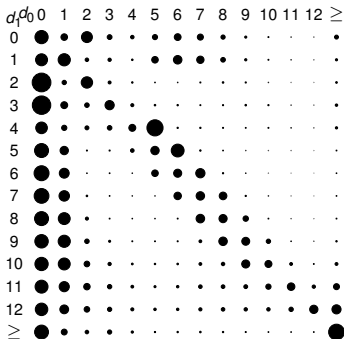
History Stack Histograms



Circle plots of stack histograms



History stack histograms consider context of occurring stack distances



⇒ increases precision of estimates



- Concrete cache set state: tuple of elements and status bits
- Two FIFO states: $[a, \underline{b}, c, d]$ and $[d, a, \underline{b}, c]$
- Only relative order w.r.t. replacement matters
- Abstract from physical positions
- ⇒ Normalized tuple: $[b, c, d, a]$
- Only ages of elements matter
- Abstract from names
- ⇒ Normalized tuple of ages, e.g. $[4, 1, 5, 0]$



- Problem: arbitrarily high ages \Rightarrow infinite model
- Introduce cutoff-age, abstract from high ages

$$\Rightarrow \alpha([4, 1, 5, 0], [6, 1, 4, 0]) = [4, 1, 4, 0]$$

- Disambiguate states by history to fit history stack histograms

\Rightarrow Abstract cache set state:

[4, 1, 4, 0]
hist: 0, 2



State space construction:

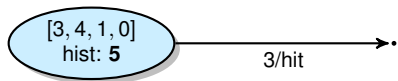
- Pick unexplored abstract state s

[3, 4, 1, 0]
hist: **5**



State space construction:

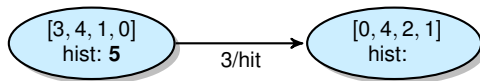
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- Enumerate all possible accesses (on ages)





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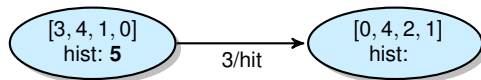
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- For each successor state s'
 - ▶ Update ages





State space construction:

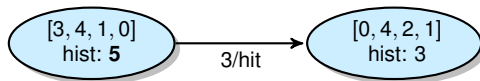
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 - ▶ Update order (apply permutation from policy table)





State space construction:

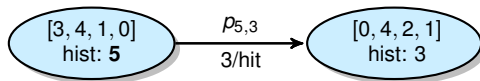
- Pick unexplored abstract state s
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 - ▶ Adjust history





State space construction:

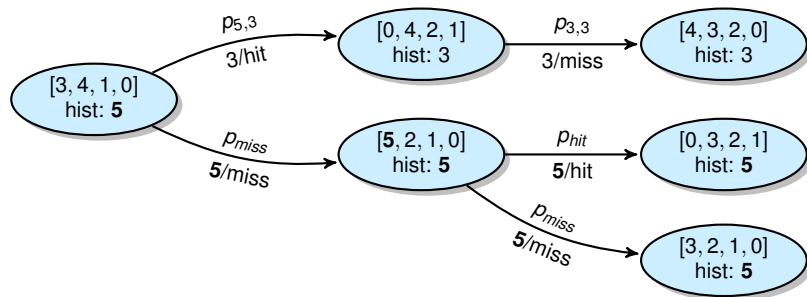
- Pick unexplored abstract state s
- Enumerate all possible accesses (on ages)
- For each successor state s'
 - ▶ Update ages
 - ▶ Update order (apply permutation from policy table)
 - ▶ Adjust history
 - ▶ Compute probability for transition $s \rightarrow s'$ from history stack histogram



Model Construction: Example



4-way FIFO, history length 1, cutoff age 5





- 1 Compute steady state probabilities $p(s_i)$ of the Markov chain
 - ▶ e.g. with the Gauss-Seidel algorithm
- 2 Compute miss probabilities $p(m_i)$
 - ▶ simple sum of values in history stack histogram

$$\Rightarrow \text{Miss ratio} = \sum_i p(s_i) \cdot p(m_i)$$



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- Shared 2nd level caches with different parameter combinations
- C programs of SPEC CPU2000
- Evaluated against full simulation runs (SimpleScalar)

- Average absolute error between 0.18% and 2.92%
- Better to spend resources on history than on higher cutoff age

- Example: 8-way associative, 1024 sets, 32 bytes line size

	LRU	PLRU	FIFO	MRU	RAND
Average error	0	0.23	0.58	2.92	1.61



- $k=4$ is handled easily
- $k=8$ often requires cheaper model parameters
i.e. stronger abstraction / less precision
- Example: Most expensive model parameters
for $k = 8$ and 2 GB main memory


	PLRU	FIFO	MRU	RAND
#states	4,166,048	8,216,745	4,129,362	2,687,856
\emptyset runtime [s]	386	440	263	234


- Comparison: \emptyset simulation time was 18 hours



- Sometimes precision is limited by resources/model parameters, e.g. 8-FIFO or 8-RAND with history
- History length is currently limited to 1.
 ≥ 2 requires $> 2\text{GB}$ main memory or stronger abstractions
- Implementation limited to policies in policy table class.
 Others (e.g. 2Q, EE-LRU, ARC) conceivable
 but most likely require different abstractions



 **Mattson, Gecsei, Slutz, and Traiger**
Evaluation techniques for storage hierarchies
IBM Systems Journal, 9(2), 1970.

 **Guo and Solihin**
An analytical model for cache replacement policy performance
SIGMETRICS Perf. Eval. Rev., 34(1), 2006.

- Mattson et al. limited to LRU, but seminal work and still useful
- Guo and Solihin
 - ▶ policies directly specified in terms of probability functions
 - ▶ hence method more efficient
 - ▶ but real-world policies (PLRU, FIFO, MRU) impossible



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- Static method for miss-ratio estimation
 - ▶ stochastic model yields precise estimates
 - ▶ allows for considering real-world policies

- History stack histograms
 - ▶ extend well-known stack histograms
 - ▶ increase precision of estimates

- Policy tables
 - ▶ provide uniform representation of policies